Towards traffic-aware routing using GPS vehicle trajectories

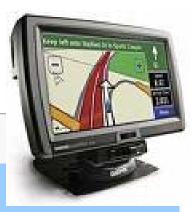
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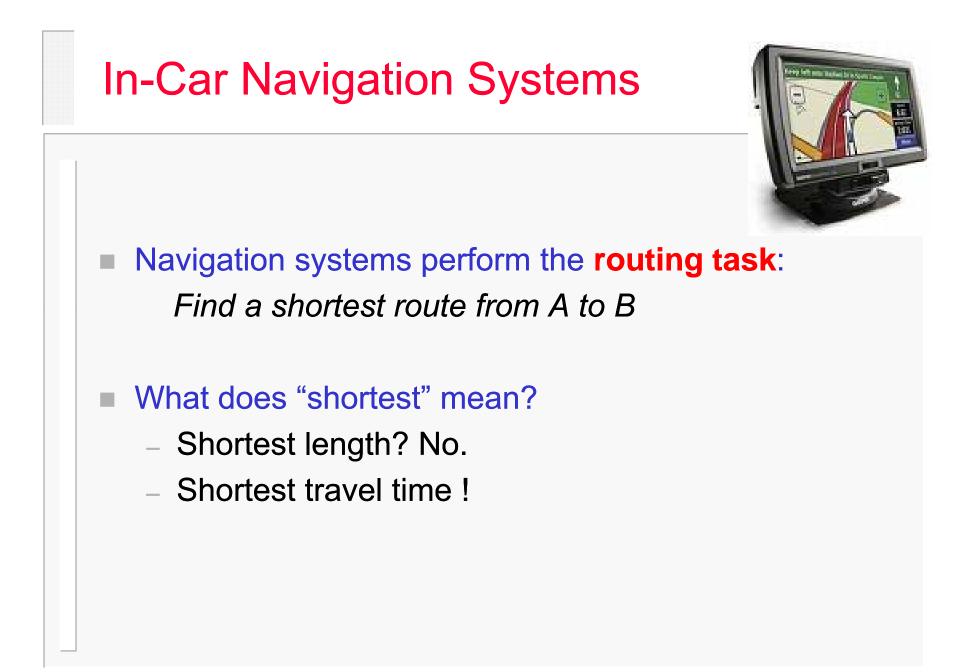




1. **Problem Description**

Enable in-car navigation systems to find the best routes using current traffic situation

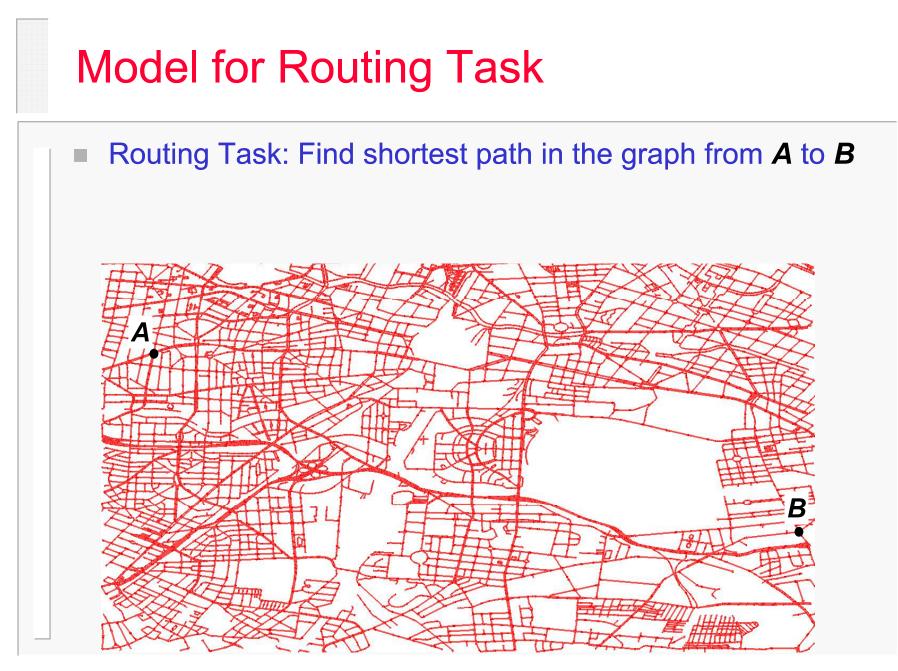
- 2. Travel Times and GPS curves
- 3. Map-Matching
 - 1. Incremental map-matching
 - 2. Global map-matching: Fréchet distance
 - 3. Global map-matching: Weak Fréchet distance
- 4. Routing System Setup and Future Work



Model for Routing Task

- Model the street network as a graph:
 - Vertices: Intersections of roads
 - Edge: A road segment between two intersections





How to Compute Shortest Paths?

- Dijkstra's shortest path algorithm:
 - Given **A**, **B**, and a graph with non-negative edge weights
 - Among all paths from A to B in the graph, compute such a path whose total weight (= sum of edge weights) is minimized
- What are our edge weights?
 - Travel times
 - The travel time on a (directed) edge from *c* to *d* is the time it takes to travel from *c* to *d*





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Travel Time of an Edge

- How do we know the travel time on an edge / road segment?
 - Usually, navigation systems derive travel times from speed limits.
 - A very smart system might take a small number of congestion points into account
 - Usually variation of travel times during rush hour is not taken into account
- New approach: Maintain a database of current travel times
 - Use GPS trajectory data from vehicle fleets (delivery trucks, taxis, etc.)
- Challenge: How do we build this database?

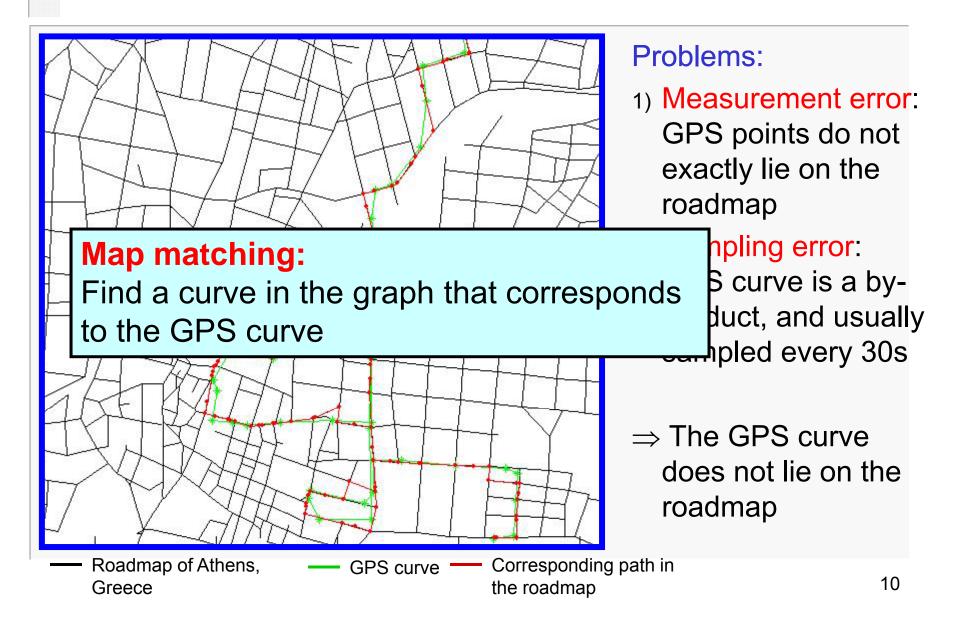
GPS Floating Car Data

Floating car data (FCD)

A sequence (trajectory) of data points, each consisting of:

- Basic vehicle telemetry, e.g., speed, direction, ABS use
- The position of the vehicle (→ tracking data) obtained by GPS tracking
- A time stamp
- Traffic assessment
 - Data from one vehicle as a sample to assess the overall traffic condition – cork swimming in the river
 - Large amounts of tracking data (e.g., taxis, public transport, utility vehicles, private vehicles)
 - \rightarrow Accurate picture of the traffic condition
 - Tracking data needs to be related to the road network
 → Map matching
 - Time stamps from FCD yield travel times for road segments

GPS Vehicle Tracking Data







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Available Map-Matching Algorithms

Incremental map-matching

- Follow greedy strategy of incrementally extending solution from an already matched edge, e.g., [BPSW05]
- No quality guarantee
- Classical approach
- Global map-matching
 - Find among all possible trajectories in the road network the one that is most similar to the vehicle trajectory
 - Distance measure assesses similarity = quality guarantee
 - Fréchet distance (strong, weak) [BPSW05, WSP06]

BPSW05: ``On Map-Matching Vehicle Tracking Data" (S. Brakatsoulas, D. Pfoser, R. Salas, and C. Wenk), *Proc. 31st Conference on Very Large Data Bases (VLDB)*: 853-864, 2005, Trondheim, Norway.

WSP06: ``Addressing the Need for Map-Matching Speed: Localizing Global Curve-Matching Algorithms"(C. Wenk and R. Salas and D. Pfoser), *Proc. 18th International Conference on Scientific and Statistical Database Management (SSDBM)*: 379-388, 2006, Vienna, Austria.





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Incremental Map-Matching

Position-by-position, edge-by-edge strategy for mapmatching

 Initialization: Find first correspondence using a spatial range query. Assume the graph has been preprocessed for spatial range queries.

 p_{i-1}

- Evaluate for each trajectory edge (or GPS point) all road network graph edges incident to the last vertex.
- Runtime O(nd + log m), with n being the number of GPS points in the trajectory, m the size of the graph, and d the maximum degree of any vertex in the graph. In practice d is a constant and O(nd) dominates the runtime.

 C_1

 d_1

d,





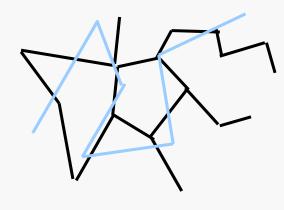
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Global Map-Matching

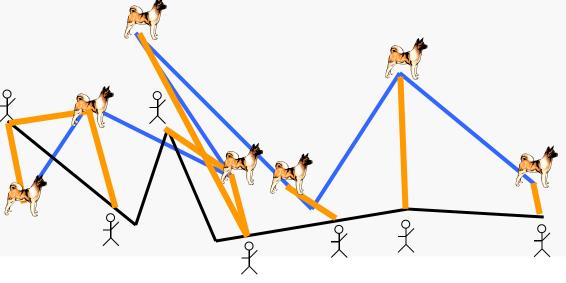
- Find a curve in the road network that is as close as possible to the vehicle trajectory
- Curves are compared using
 - Fréchet distance and
 - Weak Fréchet distance
- Minimize over all possible curves in the road network



Fréchet Distance

Dog walking example

- Person is walking his dog (person on one curve and the dog on other)
- Allowed to control their speeds but not allowed to go backwards
- Fréchet distance of the curves: *minimal leash length* necessary for both to walk the curves from beginning to end



Fréchet Distance

Fréchet Distance

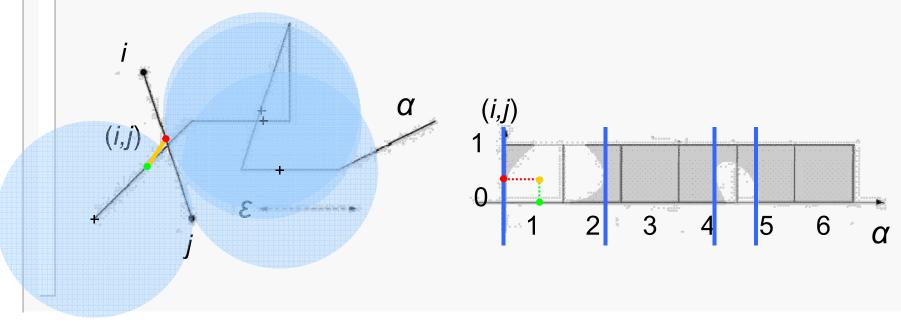
$$- \delta_{\mathsf{F}}(f,g) \coloneqq \inf_{\alpha,\beta:[0,1]\to[0,1]} \max_{t\in[0,1]} \left\| f(\alpha(t)) - g(\beta(t)) \right\|$$

- where α and β range over continuous non-decreasing reparametrizations only
- Weak Fréchet Distance
 - $\delta_F(f,g)$
 - drop the non-decreasing requirement for α and β
 - $\delta_{\mathsf{F}}(f,g) \leq \delta_{\mathsf{F}}(f,g)$
- Well-suited for the comparison of trajectories since they take the continuity of the curves into account

Free Space Diagram

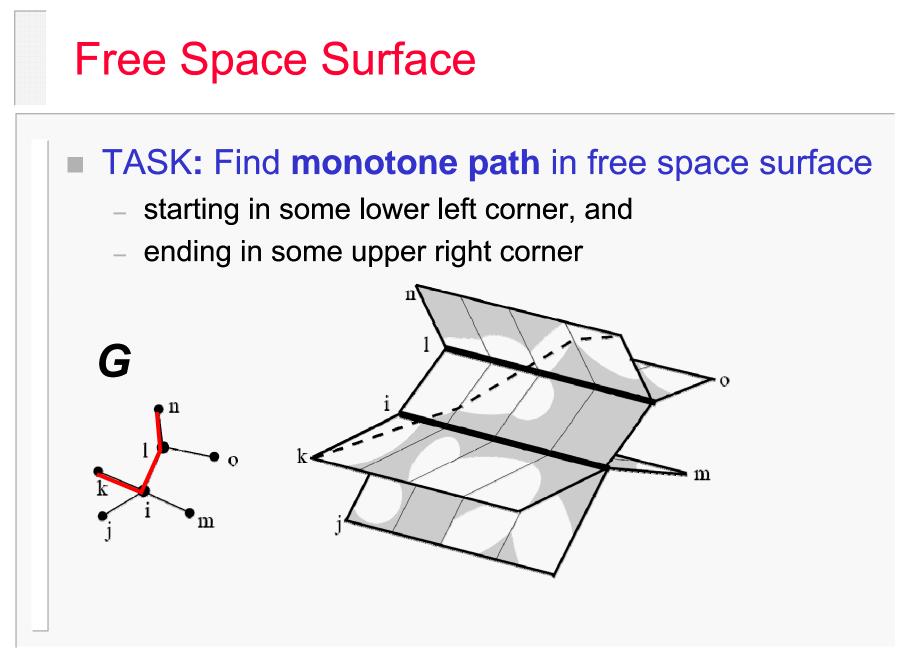
Decision variant of the global map-matching problem

- For a fixed $\varepsilon > 0$ decide whether there exists a path in the road network with distance at most ε to the vehicle trajectory α
- For each (straight-line) edge (*i*,*j*) in a graph G let its corresponding Freespace Diagram FD_{*i*,*j*} = FD(α, (*i*,*j*))



Free Space Surface

- Glue free space diagrams FD_{i,j} together according to adjacency information in the graph G
- Free space surface of trajectory α and the graph *G*



Map-Matching Using the Fréchet Distance

- Find monotone path in free space surface using dynamic programming
 - O(mn log mn) time for fixed ε, where m is the size of the graph and n the size of the trajectory
- Solve minimization problem for ε using binary search
 - O(*mn* log *mn* log *b*) time, where *b* is the desired bit precision
- Conclusion
 - Finds a curve in the graph together with a given quality guarantee (= Fréchet distance to the GPS curve)
 - The algorithm needs O(*mn*) space which for large graphs will cause a large overhead
 - Accurate map-matching results
 - Slow algorithm





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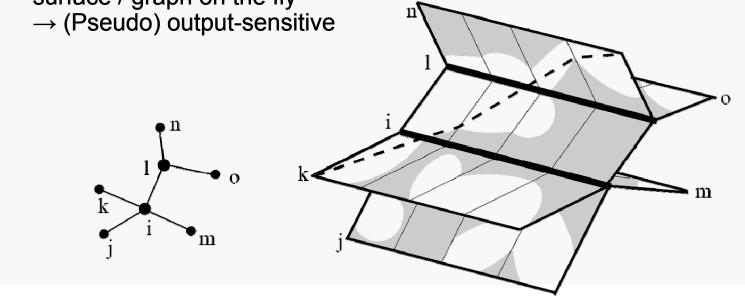
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Adaptive Clipping Map-Matching Algorithm

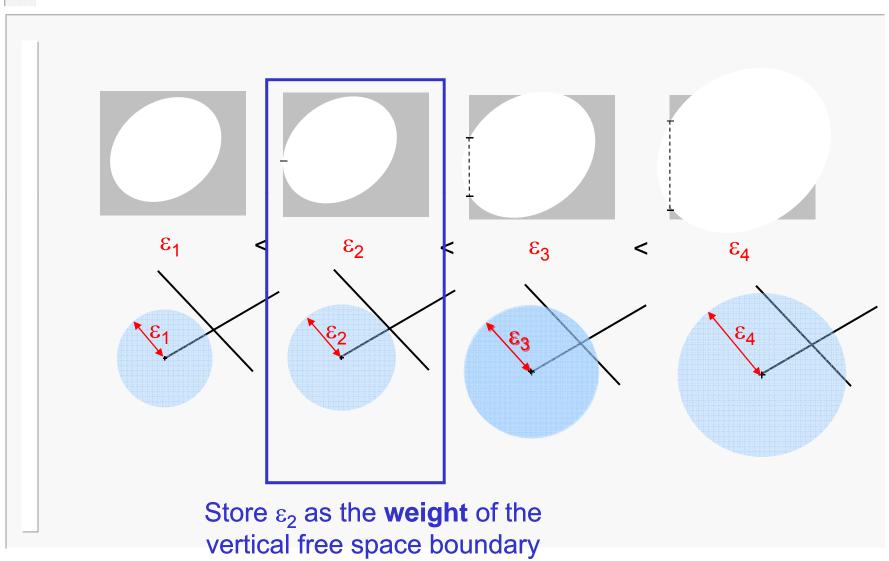
- Uses the weak Fréchet distance
- Output sensitive Map-Matching (algorithmic improvement)
 - Construct and traverse free space graph (which has quadratic complexity) on the fly
 - Improved runtime: O(K log K), where K is the size of the traversed free space which in our experiments is much smaller than *mn*
- Error-Aware Map-Matching (metadata information)
 - Use known GPS error sources to define those trajectories in the roadmap that could have led to the observed vehicle trajectory
- Adaptive Clipping algorithm:
 - Solves this error-aware map-matching task
 - Corresponds to pruning/clipping Fréchet-based algorithms

Output-Sensitive Map Matching

- The weak Fréchet distance requires finding any (possibly nonmonotone) path in the free space surface
- Express the problem as finding a shortest path in the Free Space Graph
 - Allows running Dijkstra's shortest path algorithm, which also finds ϵ (no binary search needed)
 - Allows constructing only the traversed portion of the free space surface / graph on the fly

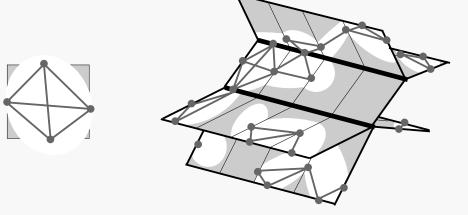


Free Space Cell



Free Space Graph

Encodes connectivity information of the free space



- For each cell boundary store optimal ε as its weight
- Weak Fréchet distance = maximum of all ε's along an optimal path in the free space graph
- \Rightarrow Need to find a shortest path (with minimum total ε)
- ⇒ Run Dijkstra's algorithm on the free space graph
- ⇒ Free space graph can be constructed on the fly

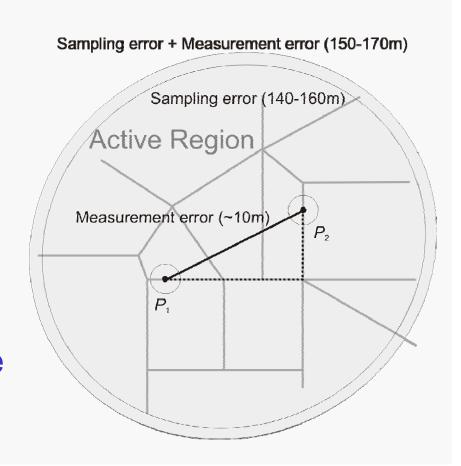
Algorithm

Running Time

- O(K log K), with K being the size of the traversed free space graph
- K = O(mn) in the worst case for traversing all the graph
- BUT, stops when shortest path to end vertex is found
- \Rightarrow Need to find a shortest path (with minimum total ε)
- Run Dijkstra's algorithm on the free space graph
- ⇒ Free space graph can be constructed on the fly

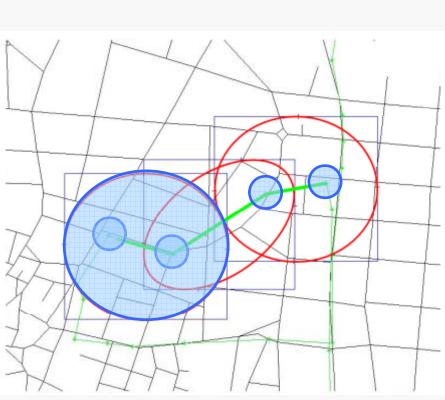
Error-Aware Map Matching

- Considering data acquisition errors
 - measurement error
 - sampling error
- Active Regions
 - areas delimiting possible positions
- Sequence of all active regions
 → error-aware representation of the vehicle trajectory



Error-Aware Map Matching

- Find curve that fulfills the following properties
 - Starts at the origin
 - Stops at the destination
 - Intersects measurement error disks around all position samples
 - Is within active regions



Adaptive Clipping Algorithm

Incremental algorithm

For each active region of an edge

- Run output-sensitive weak-Frechet/Dijkstra algorithm
- Stitch Dijkstra graphs together at active regions of vertices
- Construct overall result by tracing back stitched-together
 Dijkstra graphs
- ⇒ Solves error-aware map-matching task
- The clipping immensely reduces K. In practice K is almost constant.
- ⇒ The algorithm is very fast, almost O(n log n) runtime

Empirical Evaluation

- GPS vehicle tracking data
 - 27 trajectories
 (~16k GPS points)
 - sampling rate 30 seconds
 - max speed 80km/h
- Road network data
 - vector map of Athens, Greece
 (40 x 40km)
- Evaluating matching quality
 - results from *incremental* vs. *global* method
 - Averaged Fréchet distance (uses sampled integral / sum instead of maximum)

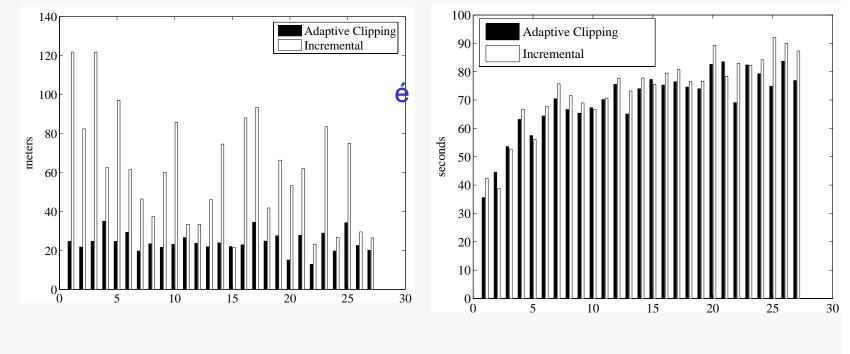
Empirical Evaluation

Matching quality

(Use average Fréchet distance as a quality measure)

Matching speed

 Adaptive Clipping runs as fast as the Incremental Algorithm



⇒ Adaptive Clipping is fast and yields good matching quality





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Routing System Setup (work in progress)

- Maintain a database of current travel times
 - Update database with current feed of GPS
 data
 - Test system in Athens gets new data every 5 minutes
 - Data provided by vehicle fleet systems which already have an uplink for GPS data transfer between the vehicle and a central server
- Communication with in-car navigation system
 - Through GPRS (cell phone) not yet implemented

central server

in-car client

Client Access to Travel Time Database (work in progress) [KW07]

- Local computation model:
 - Client computes routes itself
 - Updated travel times (in neighborhood of initial route) are transferred when they differ too much from original travel time
- Central computation model:
 - Central server computes routes

KW07: ``Dynamic Routing'' (N. Kalinowski and C. Wenk), *Technical Report CS-TR-2007-005, Department of Computer Science,* University of Texas at San Antonio, 2007.

Real-World Applications

- This is an extremely hot research area
- Almost every car manufacturer and GPS receiver manufacturer is building a *similar* prototype system
 - Different models for GPS data collection and clientserver setup / communication
 - GPS manufacturers download "old" GPS data offline from GPS receivers
 - \rightarrow Not current travel times

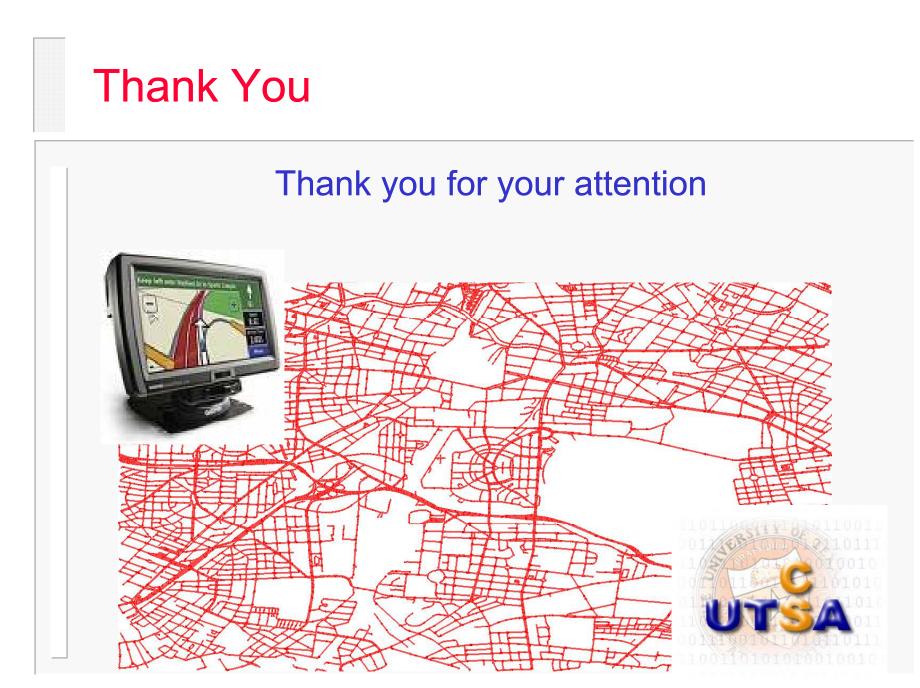
Other Types of Tracking Data

GPS drawbacks

GPS receiver needs line of sight to satellites.
 → Does not work inside buildings, and has problems in cities with very tall buildings

Other types of tracking data

- Other positioning technology (wireless networks, GSM)
 - \rightarrow Measurement error is much higher
 - \rightarrow Pilot project in Athens to collect GSM signal strength data
- Type of moving objects (planes, people)



Quality of Matching Result

- Comparing Fréchet distance of original and matched trajectory
- Fréchet distances strongly affected by outliers, since it takes the *maximum* over a set of distances.

 $\delta_{\mathsf{F}}(f,g) \coloneqq \inf_{\alpha,\beta:[0,1]\to[0,1]} \max_{t\in[0,1]} \left\| f(\alpha(t)) - g(\beta(t)) \right\|$

• Average Fréchet distance – replace the maximum with a path integral over the reparametrization curve ($\alpha(t),\beta(t)$):

$$\delta_{\mathcal{F}}(f,g) := \inf_{\alpha,\beta:[0,1]\to[0,1]} \int_{(\alpha,\beta)} \left\| f(\alpha(t)) - g(\beta(t)) \right\|$$

 Remark: Dividing by the arclength of the reparametrization curve yields a normalization, and hence an "average" of all distances.