## Towards traffic-aware routing using GPS vehicle trajectories

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## Outline

1. Problem Description

Enable in-car navigation systems to find the best routes using current traffic situation
2. Travel Times and GPS curves
3. Map-Matching

1. Incremental map-matching
2. Global map-matching: Fréchet distance
3. Global map-matching: Weak Fréchet distance
4. Routing System Setup and Future Work

## In-Car Navigation Systems

- Navigation systems perform the routing task:

Find a shortest route from $A$ to $B$

- What does "shortest" mean?
- Shortest length? No.
- Shortest travel time !


## Model for Routing Task

- Model the street network as a graph:
- Vertices: Intersections of roads
- Edge: A road segment between two intersections



## Model for Routing Task

- Routing Task: Find shortest path in the graph from $\boldsymbol{A}$ to $\boldsymbol{B}$



## How to Compute Shortest Paths?

- Dijkstra's shortest path algorithm:
- Given $\boldsymbol{A}, \boldsymbol{B}$, and a graph with non-negative edge weights
- Among all paths from $\boldsymbol{A}$ to $\boldsymbol{B}$ in the graph, compute such a path whose total weight (= sum of edge weights) is minimized
- What are our edge weights?
- Travel times
- The travel time on a (directed) edge from $\boldsymbol{c}$ to $\boldsymbol{d}$ is the time it takes to travel from $\boldsymbol{c}$ to $\boldsymbol{d}$


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## Travel Time of an Edge

- How do we know the travel time on an edge / road segment?
- Usually, navigation systems derive travel times from speed limits.
- A very smart system might take a small number of congestion points into account
- Usually variation of travel times during rush hour is not taken into account
- New approach: Maintain a database of current travel times
- Use GPS trajectory data from vehicle fleets (delivery trucks, taxis, etc.)
- Challenge: How do we build this database?


## GPS Floating Car Data

- Floating car data (FCD)

A sequence (trajectory) of data points, each consisting of:

- Basic vehicle telemetry, e.g., speed, direction, ABS use
- The position of the vehicle ( $\rightarrow$ tracking data) obtained by GPS tracking
- A time stamp
- Traffic assessment
- Data from one vehicle as a sample to assess the overall traffic condition - cork swimming in the river
- Large amounts of tracking data (e.g., taxis, public transport, utility vehicles, private vehicles)
$\rightarrow$ Accurate picture of the traffic condition
- Tracking data needs to be related to the road network $\rightarrow$ Map matching
- Time stamps from FCD yield travel times for road segments ${ }_{9}$


## GPS Vehicle Tracking Data



## Problems:

1) Measurement error: GPS points do not exactly lie on the roadmap

## Map matching:

ppling error:
Find a curve in the graph that corresponds to the GPS curve

5 curve is a byduct, and usually hpled every 30s
$\Rightarrow$ The GPS curve does not lie on the roadmap

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## Available Map-Matching Algorithms

- Incremental map-matching
- Follow greedy strategy of incrementally extending solution from an already matched edge, e.g., [BPSW05]
- No quality guarantee
- Classical approach
- Global map-matching
- Find among all possible trajectories in the road network the one that is most similar to the vehicle trajectory
- Distance measure assesses similarity = quality guarantee
- Fréchet distance (strong, weak) [BPSW05, WSP06]

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## Incremental Map-Matching

- Position-by-position, edge-by-edge strategy for mapmatching

- Initialization: Find first correspondence using a spatial range query. Assume the graph has been preprocessed for spatial range queries.
- Evaluate for each trajectory edge (or GPS point) all road network graph edges incident to the last vertex.
- Runtime $\mathrm{O}(n d+\log m)$, with $n$ being the number of GPS points in the trajectory, $m$ the size of the graph, and $d$ the maximum degree of any vertex in the graph. In practice $d$ is a constant and $\mathrm{O}(n d)$ dominates the runtime.


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## Global Map-Matching

- Find a curve in the road network that is as close as possible to the vehicle trajectory
- Curves are compared using
- Fréchet distance and
- Weak Fréchet distance
- Minimize over all possible curves in the road network



## Fréchet Distance

- Dog walking example
- Person is walking his dog (person on one curve and the dog on other)
- Allowed to control their speeds but not allowed to go backwards
- Fréchet distance of the curves: minimal /leash length necessary for both to walk the curves from beginning to end



## Fréchet Distance

- Fréchet Distance
- $\delta_{F}(f, g):=\inf _{\alpha, \beta: 0,1] \rightarrow[0,1]} \max _{t \in[0,1]}\|f(\alpha(t))-g(\beta(t))\|$
- where $\alpha$ and $\beta$ range over continuous non-decreasing reparametrizations only
- Weak Fréchet Distance
- $\tilde{\delta}_{F}(f, g)$
- drop the non-decreasing requirement for $\alpha$ and $\beta$
${ }^{-} \delta_{F}(f, g) \leq \delta_{F}(f, g)$
- Well-suited for the comparison of trajectories since they take the continuity of the curves into account


## Free Space Diagram

Decision variant of the global map-matching problem

- For a fixed $\varepsilon>0$ decide whether there exists a path in the road network with distance at most $\varepsilon$ to the vehicle trajectory $\alpha$
- For each (straight-line) edge ( $i, j$ ) in a graph $G$ let its corresponding Freespace Diagram $\mathrm{FD}_{i, j}=\mathrm{FD}(\alpha,(i, j))$



## Free Space Surface

- Glue free space diagrams $\mathrm{FD}_{i, j}$ together according to adjacency information in the graph $G$
- Free space surface of trajectory $\alpha$ and the graph G

G


## Free Space Surface

- TASK: Find monotone path in free space surface
- starting in some lower left corner, and
- ending in some upper right corner


## G



## Map-Matching Using the Fréchet Distance

- Find monotone path in free space surface using dynamic programming
- $\mathrm{O}(m n \log m n)$ time for fixed $\varepsilon$, where $m$ is the size of the graph and $n$ the size of the trajectory
- Solve minimization problem for $\varepsilon$ using binary search
- O(mn $\log m n \log b)$ time, where $b$ is the desired bit precision
- Conclusion
- Finds a curve in the graph together with a given quality guarantee (= Fréchet distance to the GPS curve)
- The algorithm needs $\mathrm{O}(m n)$ space which for large graphs will cause a large overhead
$\Rightarrow$ Accurate map-matching results
$\Rightarrow$ Slow algorithm


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## Adaptive Clipping Map-Matching Algorithm

- Uses the weak Fréchet distance
- Output sensitive Map-Matching (algorithmic improvement)
- Construct and traverse free space graph (which has quadratic complexity) on the fly
- Improved runtime: $\mathrm{O}(\mathrm{K} \log \mathrm{K})$, where K is the size of the traversed free space which in our experiments is much smaller than $m n$
- Error-Aware Map-Matching (metadata information)
- Use known GPS error sources to define those trajectories in the roadmap that could have led to the observed vehicle trajectory
- Adaptive Clipping algorithm:
- Solves this error-aware map-matching task
- Corresponds to pruning/clipping Fréchet-based algorithms


## Output-Sensitive Map Matching

- The weak Fréchet distance requires finding any (possibly nonmonotone) path in the free space surface
- Express the problem as finding a shortest path in the Free Space Graph
- Allows running Dijkstra's shortest path algorithm, which also finds $\varepsilon$ (no binary search needed)
- Allows constructing only the traversed portion of the free space surface / graph on the fly $\rightarrow$ (Pseudo) output-sensitive



## Free Space Cell



Store $\varepsilon_{2}$ as the weight of the vertical free space boundary

## Free Space Graph

- Encodes connectivity information of the free space

- For each cell boundary store optimal $\varepsilon$ as its weight
- Weak Fréchet distance = maximum of all $\varepsilon$ 's along an optimal path in the free space graph
$\Rightarrow$ Need to find a shortest path (with minimum total $\varepsilon$ )
$\Rightarrow$ Run Dijkstra's algorithm on the free space graph
$\Rightarrow$ Free space graph can be constructed on the fly


## Algorithm

- Running Time
- $\mathrm{O}(K \log K)$, with $K$ being the size of the traversed free space graph
- $\mathrm{K}=\mathrm{O}(m n)$ in the worst case for traversing all the graph
- BUT, stops when shortest path to end vertex is found
$\Rightarrow$ Need to find a shortest path (with minimum total $\varepsilon$ )
$\Rightarrow$ Run Dijkstra's algorithm on the free space graph
$\Rightarrow$ Free space graph can be constructed on the fly


## Error-Aware Map Matching

- Considering data acquisition errors
- measurement error
- sampling error
- Active Regions
- areas delimiting possible positions
- Sequence of all active regions
$\rightarrow$ error-aware representation of the vehicle trajectory

Sampling error + Measurement error ( $150-170 \mathrm{~m}$ )


## Error-Aware Map Matching

- Find curve that fulfills the following properties
- Starts at the origin
- Stops at the destination
- Intersects measurement error disks around all position samples
- Is within active regions



## Adaptive Clipping Algorithm

- Incremental algorithm

For each active region of an edge

- Run output-sensitive weak-Frechet/Dijkstra algorithm
- Stitch Dijkstra graphs together at active regions of vertices
- Construct overall result by tracing back stitched-together Dijkstra graphs
$\Rightarrow$ Solves error-aware map-matching task
$\Rightarrow$ The clipping immensely reduces K . In practice K is almost constant.
$\Rightarrow$ The algorithm is very fast, almost $\mathrm{O}(\mathrm{n} \log \mathrm{n})$ runtime


## Empirical Evaluation

- GPS vehicle tracking data
- 27 trajectories (~16k GPS points)
- sampling rate 30 seconds
- max speed 80km/h
- Road network data
- vector map of Athens, Greece (40 x 40km)
- Evaluating matching quality
- results from incremental vs. global method
- Averaged Fréchet distance (uses sampled integral / sum instead of maximum)


## Empirical Evaluation

- Matching quality
(Use average Fréchet distance as a quality measure)

- Matching speed
- Adaptive Clipping runs as fast as the Incremental Algorithm

$\Rightarrow$ Adaptive Clipping is fast and yields good matching quality


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## Routing System Setup (work in progress)

- Maintain a database of current travel times
- Update database with current feed of GPS data
- Test system in Athens gets new data every 5 minutes
- Data provided by vehicle fleet systems which already have an uplink for GPS data transfer between the vehicle and a central server
- Communication with in-car navigation system
- Through GPRS (cell phone) - not yet implemented
central
server
in-car client


## Client Access to Travel Time Database (work in progress) [KW07]

- Local computation model:
- Client computes routes itself
- Updated travel times (in neighborhood of initial route) are transferred when they differ too much from original travel time
- Central computation model:
- Central server computes routes

KW07: `Dynamic Routing" (N. Kalinowski and C. Wenk), Technical Report CS-TR-2007-005, Department of Computer Science, University of Texas at San Antonio, 2007.

## Real-World Applications

- This is an extremely hot research area
- Almost every car manufacturer and GPS receiver manufacturer is building a similar prototype system
- Different models for GPS data collection and clientserver setup / communication
- GPS manufacturers download "old" GPS data offline from GPS receivers
$\rightarrow$ Not current travel times


## Other Types of Tracking Data

## - GPS drawbacks

- GPS receiver needs line of sight to satellites. $\rightarrow$ Does not work inside buildings, and has problems in cities with very tall buildings
- Other types of tracking data
- Other positioning technology (wireless networks, GSM) $\rightarrow$ Measurement error is much higher
$\rightarrow$ Pilot project in Athens to collect GSM signal strength data
- Type of moving objects (planes, people)


## Thank You

## Thank you for your attention



## Quality of Matching Result

- Comparing Fréchet distance of original and matched trajectory
- Fréchet distances strongly affected by outliers, since it takes the maximum over a set of distances.
$\delta_{F}(f, g):=\inf _{\alpha, \beta:[0,1] \rightarrow[0,1]} \max _{t \in[0,1]}\|f(\alpha(t))-g(\beta(t))\|$
- Average Fréchet distance - replace the maximum with a path integral over the reparametrization curve $(\alpha(\mathrm{t}), \beta(\mathrm{t}))$ :
$\delta_{F}(f, g):=\inf _{\alpha, \beta: 0,1] \rightarrow[0,1]} \int_{(\alpha, \beta)}\|f(\alpha(t))-g(\beta(t))\|$
- Remark: Dividing by the arclength of the reparametrization curve yields a normalization, and hence an „average" of all distances.


[^0]:    BPSW05: ‘`On Map-Matching Vehicle Tracking Data" (S. Brakatsoulas, D. Pfoser, R. Salas, and C. Wenk), Proc. 31st Conference on Very Large Data Bases (VLDB): 853-864, 2005, Trondheim, Norway.

    WSP06: `’Addressing the Need for Map-Matching Speed: Localizing Global Curve-Matching Algorithms"(C. Wenk and R. Salas and D. Pfoser), Proc. 18th International Conference on Scientific and Statistical Database
    Management (SSDBM): 379-388, 2006, Vienna, Austria.

