

8. Homework

Due **Tuesday 4/13/10** before class

1. Kruskal's tree (4 points)

Since Kruskal's algorithm computes a forest of trees which are incrementally merged, it is not immediately obvious why the output is actually a tree.

Argue why Kruskal's algorithm computes a tree (= connected acyclic graph). Do not use the fact that somebody else has proven that Kruskal's algorithm computes an MST.

2. Negative edge weights (2+2+4 points)

- a) Give an example of a directed connected graph with real edge weights (that may be negative) for which Dijkstra's algorithm produces incorrect answers. Justify your answer.
- b) Does Dijkstra's algorithm only produce incorrect answers in the presence of a negative weight cycle, or could it also produce incorrect answers in the mere presence of negative weight edges (without any negative weight cycles)? Justify your answer.
- c) Suppose the weighted, directed graph $G = (V, E)$ has a special structure in which edges that leave the source vertex s may have negative weights. All other edge weights are nonnegative, and there are no negative-weight cycles. Show that Dijkstra's algorithm correctly finds shortest paths from s in G .

3. Adding an edge (4 points)

Let $G = (V, E)$ be a connected undirected graph with weight function $w : E \rightarrow \mathbb{R}_0^+$ (i.e., all edge weights are ≥ 0). Assume edge weights are distinct. Further, let a minimum spanning tree T on G be given.

Now, assume that one new edge (u, v) , with $u, v \in V$, with weight $w(u, v)$ is added to G . (This weight is different from all other edge weights.)

Give an efficient algorithm to test if T remains the minimum spanning tree for this new graph. Your algorithm should run in $O(|E|)$ time. Can you make it run in $O(|V|)$ time?

4. Faster MST (8 points)

Let $G = (V, E)$ be a connected undirected graph with edge weights $w : E \rightarrow \mathbb{R}$.

- (a) If all of the edge weights are integers between 1 and $|V|$, how fast can the minimum spanning tree be computed? (Give the *most efficient* algorithm you can think of.)
- (b) Now assume all edge weights are integers between 1 and 10. Show that Prim's algorithm can be implemented to work in $O(|V| + |E|)$ time in this case. (*Hint: Suggest a data structure that replaces the priority queue and analyze the runtime.*)