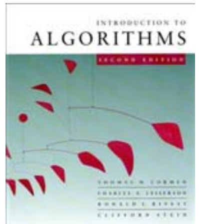
CS 5633 -- Spring 2009



Range Trees

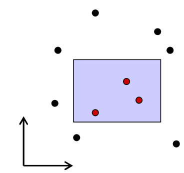
Carola Wenk

Slides courtesy of Charles Leiserson with small changes by Carola Wenk

Orthogonal range searching

Input: n points in d dimensions
- E.g., representing a database of n records each with d numeric fields

Query: Axis-aligned ***box*** (in 2D, a rectangle)
- Report on the points inside the box:
  - Are there any points?
  - How many are there?
  - List the points.

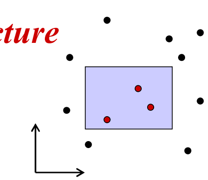


Orthogonal range searching

Input: n points in d dimensions

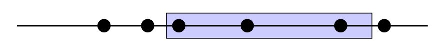
Query: Axis-aligned ***box*** (in 2D, a rectangle)
- Report on the points inside the box

Goal: Preprocess points into a data structure to support fast queries
- Primary goal: ***Static data structure***
- In 1D, we will also obtain a dynamic data structure supporting insert and delete



1D range searching

In 1D, the query is an interval:

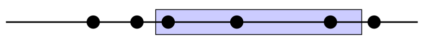


First solution:
- Sort the points and store them in an array
  - Solve query by binary search on endpoints.
  - Obtain a static structure that can list k answers in a query in $O(k + \log n)$ time.

Goal: Obtain a dynamic structure that can list k answers in a query in $O(k + \log n)$ time.

1D range searching

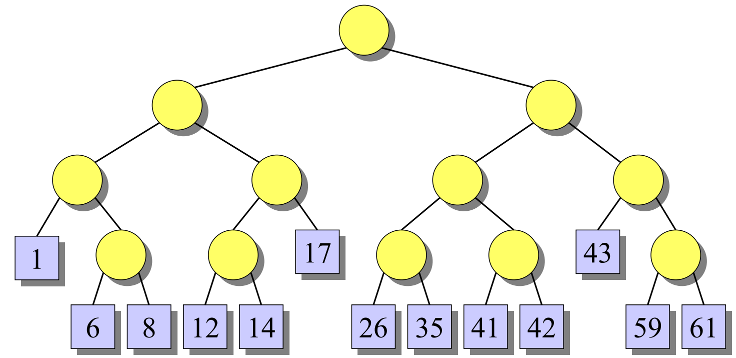
In 1D, the query is an interval:



New solution that extends to higher dimensions:

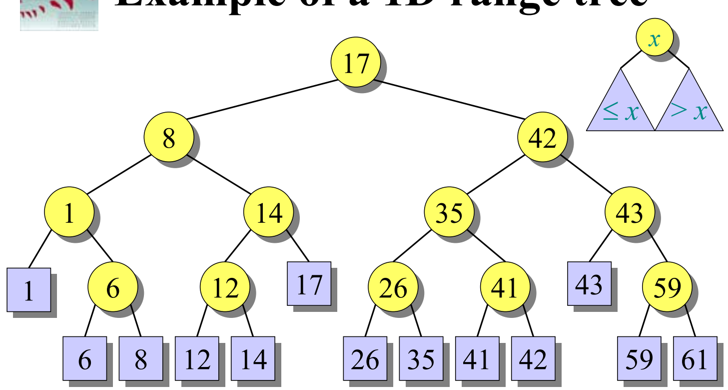
- Balanced binary search tree
  - New organization principle:
    Store points in the *leaves* of the tree.
  - Internal nodes store copies of the leaves to satisfy binary search property:
    - Node *x* stores in *key*[*x*] the maximum key of any leaf in the left subtree of *x*.

Example of a 1D range tree



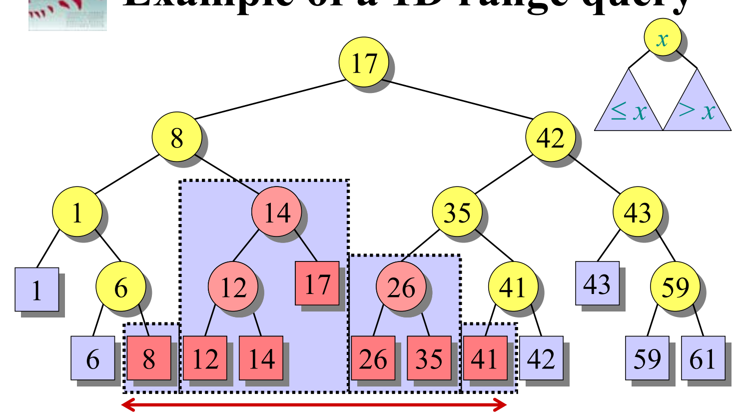
key[*x*] is the maximum key of any leaf in the left subtree of *x*.

Example of a 1D range tree



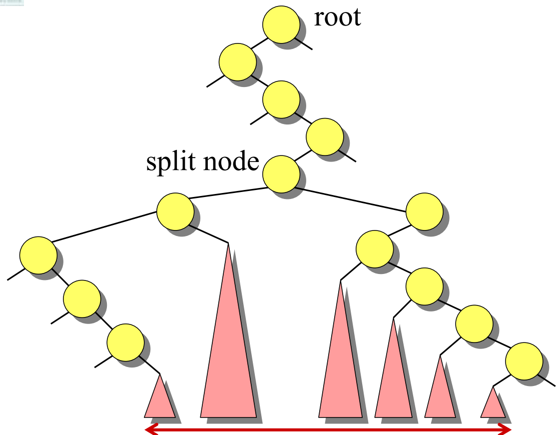
key[*x*] is the maximum key of any leaf in the left subtree of *x*.

Example of a 1D range query



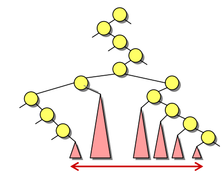
RANGE-QUERY([7, 41])

General 1D range query



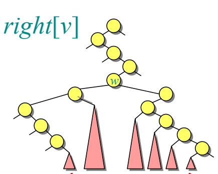
Pseudocode, part 1: Find the split node

1D-RANGE-QUERY(T, $[x_1, x_2]$)
  $w \leftarrow$ root[T]
  while w is not a leaf and ($x_2 \leq key[w]$ or $key[w] < x_1$)
    do if $x_2 \leq key[w]$
      then $w \leftarrow left[w]$
      else $w \leftarrow right[w]$
  // w is now the split node
  [*traverse left and right from w and report relevant subtrees*]



Pseudocode, part 2: Traverse left and right from split node

1D-RANGE-QUERY(T, $[x_1, x_2]$)
  [*find the split node*]
  // w is now the split node
  if w is a leaf
  then output the leaf w if $x_1 \leq key[w] \leq x_2$
  else $v \leftarrow left[w]$ // Left traversal
    while v is not a leaf
      do if $x_1 \leq key[v]$
        then output the subtree rooted at $right[v]$
          $v \leftarrow left[v]$
        else $v \leftarrow right[v]$
    output the leaf v if $x_1 \leq key[v] \leq x_2$
    [*symmetrically for right traversal*]

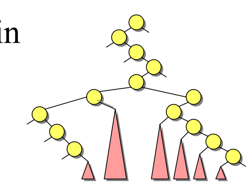


Analysis of 1D-RANGE-QUERY

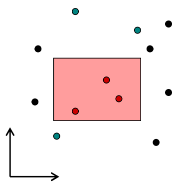
Query time: Answer to range query represented by O(log n) subtrees found in O(log n) time. Thus:

- Can test for points in interval in O(log n) time.
- Can report all k points in interval in O(k + log n) time.
- Can count points in interval in O(log n) time

Space: O(n)
Preprocessing time: O(n log n)



2D range trees

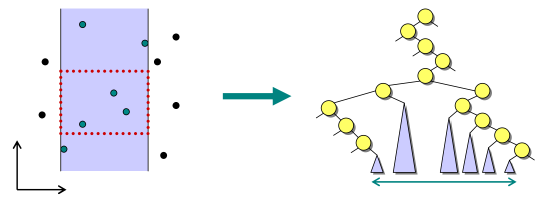


2D range trees

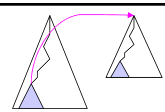
Store a ***primary*** 1D range tree for all the points based on *x*-coordinate.

Thus in O(log *n*) time we can find O(log *n*) subtrees representing the points with proper *x*-coordinate.

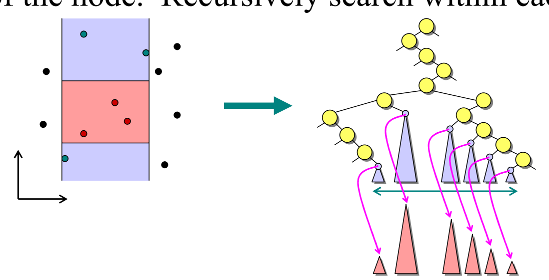
How to restrict to points with proper *y*-coordinate?



2D range trees

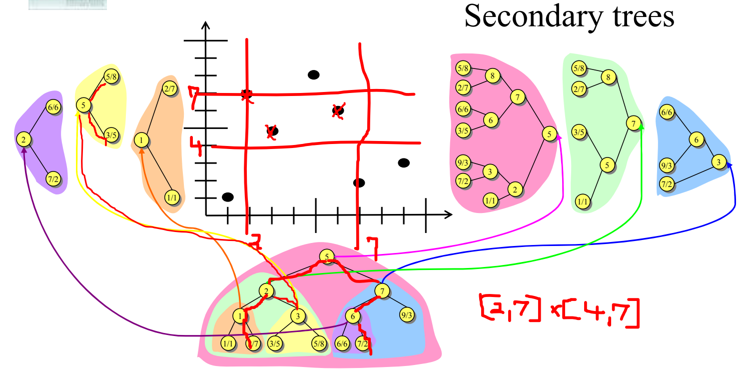


Idea: In primary 1D range tree of *x*-coordinate, **<u>every</u>** node stores a ***secondary*** 1D range tree based on *y*-coordinate for all points in the subtree of the node. Recursively search within each.



2D range tree example

Secondary trees



Primary tree

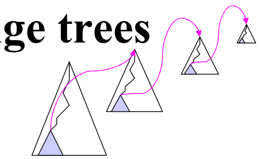
Analysis of 2D range trees

Query time: In $O(\log^2 n) = O((\log n)^2)$ time, we can represent answer to range query by $O(\log^2 n)$ subtrees. Total cost for reporting k points: $O(k + (\log n)^2)$.

Space: The secondary trees at each level of the primary tree together store a copy of the points. Also, each point is present in each secondary tree along the path from the leaf to the root. Either way, we obtain that the space is $O(n \log n)$.

Preprocessing time: $O(n \log n)$

d-dimensional range trees



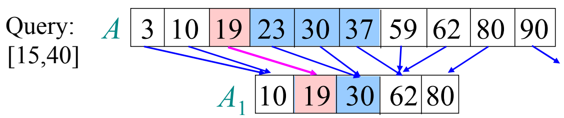
Each node of the secondary y-structure stores a tertiary z-structure representing the points in the subtree rooted at the node, etc.

Save one log factor using fractional cascading

Query time: $O(k + \log^d n)$ to report k points.
Space: $O(n \log^{d-1} n)$
Preprocessing time: $O(n \log^{d-1} n)$

Search in Subsets

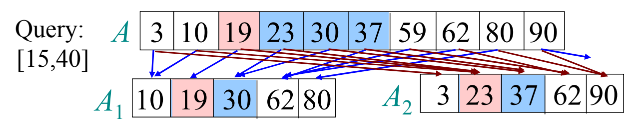
Given: Two sorted arrays A_1 and A, with $A_1 \subseteq A$
A query interval $[l,r]$
Task: Report all elements e in A_1 and A with $l \leq e \leq r$
Idea: Add pointers from A to A_1:
$\rightarrow$ For each $a \in A$ add a pointer to the smallest element $b \in A_1$ with $b \geq a$
Query: Find $l \in A$, follow pointer to A_1. Both in A and A_1 sequentially output all elements in $[l,r]$.



Runtime: $O((\log n + k) + (1 + k)) = O(\log n + k)$

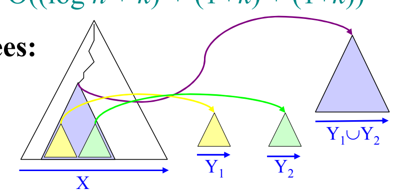
Search in Subsets (cont.)

Given: Three sorted arrays A_1, A_2, and A, with $A_1 \subseteq A$ and $A_2 \subseteq A$



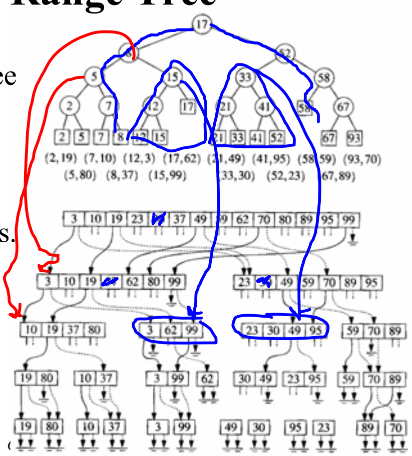
Runtime: $O((\log n + k) + (1+k) + (1+k)) = O(\log n + k)$

Range trees:



Fractional Cascading: Layered Range Tree

Replace 2D range tree with a layered range tree, using sorted arrays and pointers instead of the secondary range trees.



Preprocessing:

$O(n \log n)$

Query:

$O(\log n + k)$

d-dimensional range trees

Query time: $O(k + \log^{d-1} n)$ to report k points, uses fractional cascading in the last dimension

Space: $O(n \log^{d-1} n)$

Preprocessing time: $O(n \log^{d-1} n)$

Best data structure to date:

Query time: $O(k + \log^{d-1} n)$ to report k points.

Space: $O(n (\log n / \log \log n)^{d-1})$

Preprocessing time: $O(n \log^{d-1} n)$