CS 3343 – Fall 2011

INTRODUCTION TO
ALGORITHMS
SECOND EDITION

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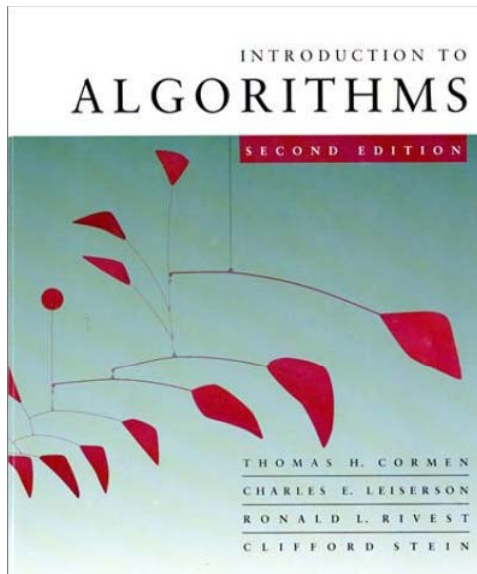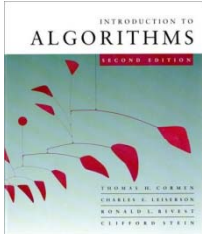CHARLES E. LEISERSON
RONALD L. RIVEST
CLIFFORD STEIN

More Divide & Conquer

Carola Wenk

Slides courtesy of Charles Leiserson with small changes by Carola Wenk
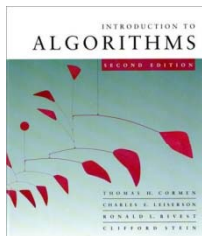
Powering a number

Problem: Compute a^n, where $n \in \mathbf{N}$.

Naive algorithm: $\Theta(n)$.

Divide-and-conquer algorithm: (recursive squaring)

$$a^n = \begin{cases} a^{n/2} \cdot a^{n/2} & \text{if } n \text{ is even;} \\ a^{(n-1)/2} \cdot a^{(n-1)/2} \cdot a & \text{if } n \text{ is odd.} \end{cases}$$

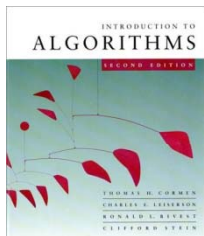$$T(n) = T(n/2) + \Theta(1) \implies T(n) = \Theta(\log n) \, .$$

Fibonacci numbers

Recursive definition:

$$F_n = \begin{cases} 0 & \text{if } n = 0; \\ 1 & \text{if } n = 1; \\ F_{n-1} + F_{n-2} & \text{if } n \geq 2. \end{cases}$$

0   1   1   2   3   5   8   13 21 34 $\cdots$

Naive recursive algorithm: $\Omega(\phi^n)$ (exponential time), where $\phi = (1 + \sqrt{5})/2$ is the *golden ratio*.
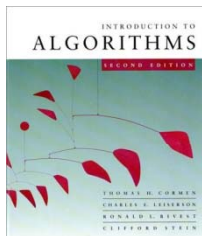
Computing Fibonacci numbers

Naive recursive squaring:

$$F_n = \phi^n / \sqrt{5} \text{ rounded to the nearest integer.}$$

- Recursive squaring: $\Theta(\log n)$ time.
- This method is unreliable, since floating-point arithmetic is prone to round-off errors.

Bottom-up (one-dimensional dynamic programming):

- Compute $F_0, F_1, F_2, \ldots, F_n$ in order, forming each number by summing the two previous.
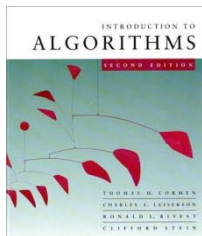- Running time: $\Theta(n)$.

Convex Hull

- Given a set of pins on a pinboard

- And a rubber band around them

- How does the rubber band look when it snaps tight?

- We represent convex hull as the sequence of points on the convex hull polygon, in counter-clockwise order.
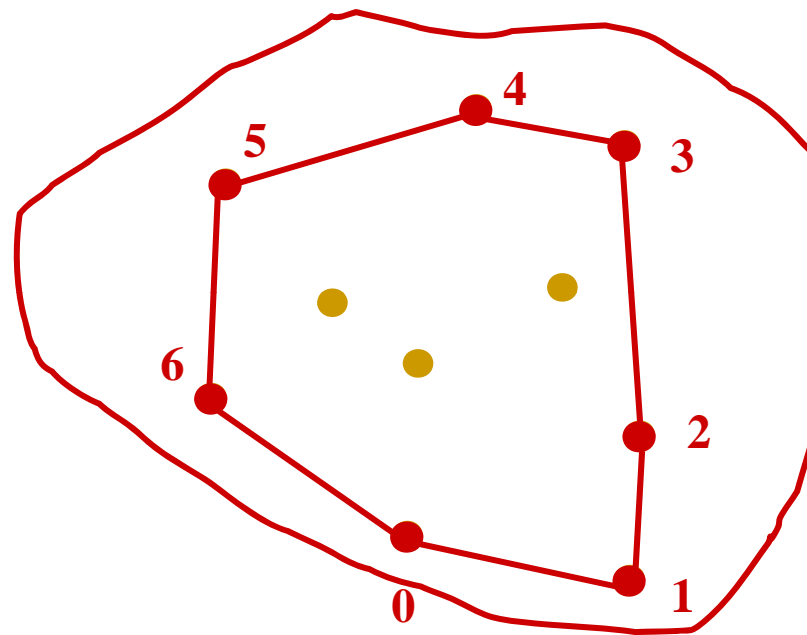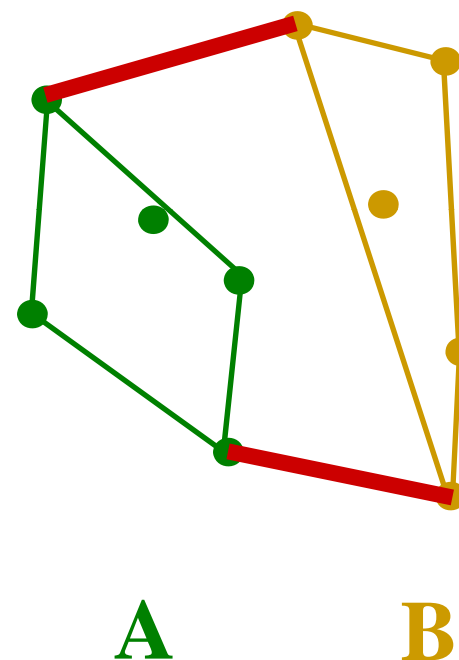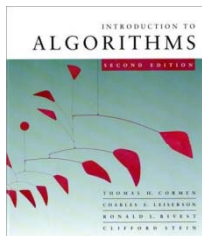
Convex Hull: Divide & Conquer

- Preprocessing: sort the points by x-coordinate

- Divide the set of points into two sets **A** and **B**:

  - **A** contains the left $\lfloor n/2 \rfloor$ points,

  - **B** contains the right $\lceil n/2 \rceil$ points

- Recursively compute the convex hull of **A**

- Recursively compute the convex hull of **B**

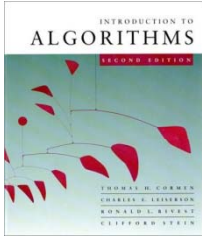- Merge the two convex hulls

A   **B**
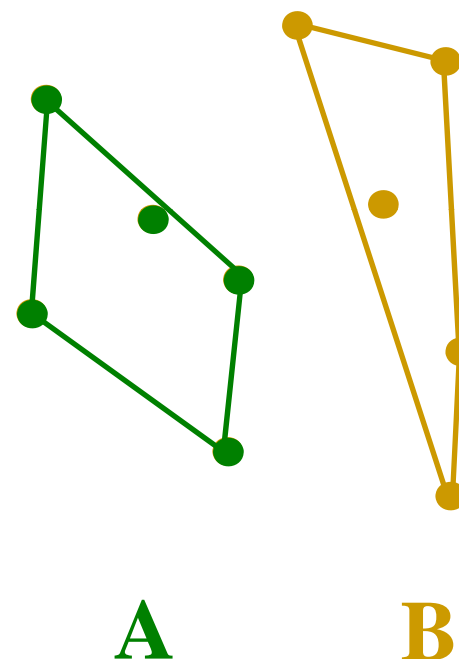
Merging

- **Find upper and lower tangent**

- With those tangents the convex hull of A∪B can be computed from the convex hulls of A and the convex hull of B in O(n) linear time

A   **B**

Finding the lower tangent

a = rightmost point of A

b = leftmost point of B

while T=ab not lower tangent to both
    convex hulls of A and B do{

    while T not lower tangent to
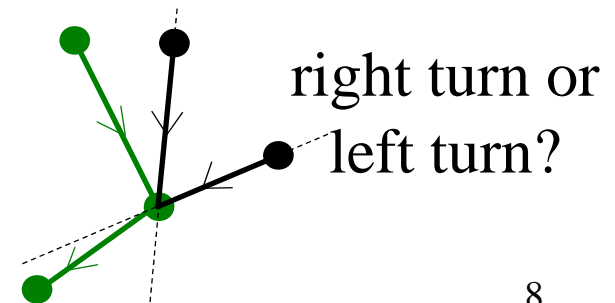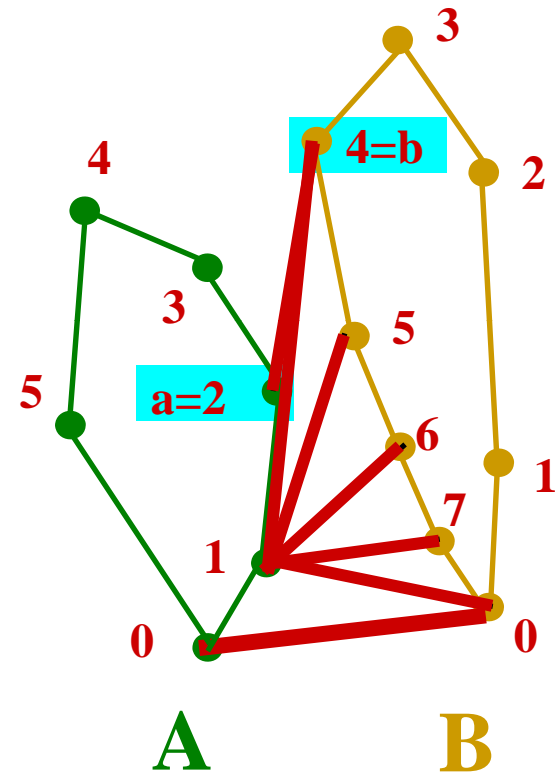   convex hull of A do{
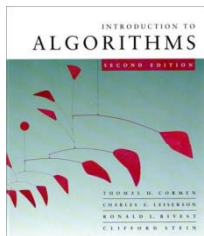   a=a-1
    }
    while T not lower tangent to
    convex hull of B do{
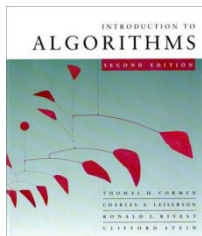   b=b+1
    }
}

can be checked
in constant time

3

4=b

4

2

3

5

5

a=2

6

1

7

1

0

0

A   **B**

right turn or
left turn?

Convex Hull: Runtime

● Preprocessing: sort the points by x-coordinate

$O(n \log n)$ just once

● Divide the set of points into two sets **A** and **B**:

$O(1)$

  ● **A** contains the left $\lfloor n/2 \rfloor$ points,

  ● **B** contains the right $\lceil n/2 \rceil$ points

●Recursively compute the convex hull of **A**

$T(n/2)$

●Recursively compute the convex hull of **B**
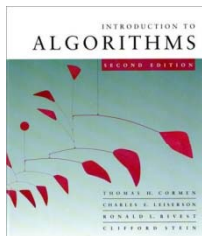
$T(n/2)$

● Merge the two convex hulls

$O(n)$

Convex Hull: Runtime

- Runtime Recurrence:

  $$T(n) = 2\ T(n/2) + dn$$

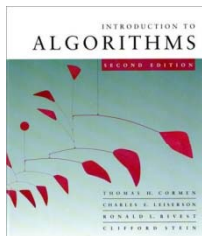- Solves to $T(n) = \Theta(n \log n)$

Matrix multiplication

Input: $A = [a_{ij}], B = [b_{ij}].$

Output: $C = [c_{ij}] = A \cdot B.$ $\quad\Big\} \quad i, j = 1, 2, \ldots, n.$

$$
\begin{bmatrix}
c_{11} & c_{12} & \cdots & c_{1n} \\
c_{21} & c_{22} & \cdots & c_{2n} \\
\vdots & \vdots & \ddots & \vdots \\
c_{n1} & c_{n2} & \cdots & c_{nn}
\end{bmatrix}
=
\begin{bmatrix}
a_{11} & a_{12} & \cdots & a_{1n} \\
a_{21} & a_{22} & \cdots & a_{2n} \\
\vdots & \vdots & \ddots & \vdots \\
a_{n1} & a_{n2} & \cdots & a_{nn}
\end{bmatrix}
\cdot
\begin{bmatrix}
b_{11} & b_{12} & \cdots & b_{1n} \\
b_{21} & b_{22} & \cdots & b_{2n} \\
\vdots & \vdots & \ddots & \vdots \\
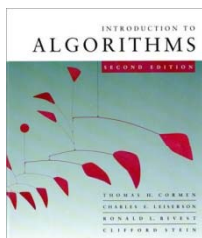b_{n1} & b_{n2} & \cdots & b_{nn}
\end{bmatrix}
$$

$$
c_{ij} = \sum_{k=1}^{n} a_{ik} \cdot b_{kj}
$$

Standard algorithm

for $i \leftarrow 1$ **to** n
    do for $j \leftarrow 1$ **to** n
        do $c_{ij} \leftarrow 0$
            for $k \leftarrow 1$ **to** n
                do $c_{ij} \leftarrow c_{ij} + a_{ik} \cdot b_{kj}$

Running time $- \Theta(n^3)$
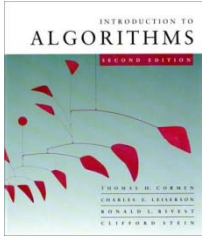
Divide-and-conquer algorithm

IDEA:

$n{\times}n$ matrix = $2{\times}2$ matrix of $(n/2){\times}(n/2)$ submatrices:

$$\begin{bmatrix} r & s \\ t & u \end{bmatrix} = \begin{bmatrix} a & b \\ c & d \end{bmatrix} \cdot \begin{bmatrix} e & f \\ g & h \end{bmatrix}$$

$$C \quad = \quad A \quad \cdot \quad B$$

$$r = a{\cdot}e + b{\cdot}g$$
$$s = a{\cdot}f + b{\cdot}h$$
$$t = c{\cdot}e + d{\cdot}g$$
$$u = c{\cdot}f + d{\cdot}h$$

8 recursive mults of $(n/2){\times}(n/2)$ submatrices

4 adds of $(n/2){\times}(n/2)$ submatrices

Analysis of D&C algorithm
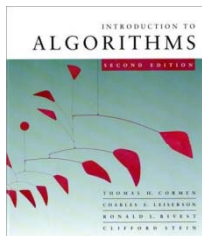
$$T(n) = 8\,T(n/2) + \Theta(n^2)$$

 # submatrices

submatrix size

work adding submatrices

$$n^{\log_b a} = n^{\log_2 8} = n^3 \;\Rightarrow\; \text{CASE 1} \;\Rightarrow\; T(n) = \Theta(n^3).$$

No better than the ordinary algorithm.

Strassen's idea

- Multiply 2×2 matrices with only 7 recursive mults.

$$P_1 = a \cdot (f - h)$$
$$P_2 = (a + b) \cdot h$$
$$P_3 = (c + d) \cdot e$$
$$P_4 = d \cdot (g - e)$$
$$P_5 = (a + d) \cdot (e + h)$$
$$P_6 = (b - d) \cdot (g + h)$$
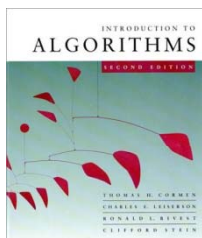$$P_7 = (a - c) \cdot (e + f)$$

$$r = P_5 + P_4 - P_2 + P_6$$
$$s = P_1 + P_2$$
$$t = P_3 + P_4$$
$$u = P_5 + P_1 - P_3 - P_7$$

7 mults, 18 adds/subs.
Note: No reliance on commutativity of mult!

Strassen's idea

- Multiply 2×2 matrices with only 7 recursive mults.

$$P_1 = a \cdot (f - h)$$
$$P_2 = (a + b) \cdot h$$
$$P_3 = (c + d) \cdot e$$
$$P_4 = d \cdot (g - e)$$
$$P_5 = (a + d) \cdot (e + h)$$
$$P_6 = (b - d) \cdot (g + h)$$
$$P_7 = (a - c) \cdot (e + f)$$

$$
\begin{aligned}
r &= P_5 + P_4 - P_2 + P_6 \\
&= (a + d)(e + h) \\
&\quad + d(g - e) - (a + b)h \\
&\quad + (b - d)(g + h) \\
&= ae + ah + de + dh \\
&\quad + dg - de - ah - bh \\
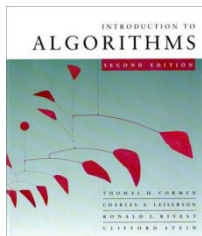&\quad + bg + bh - dg - dh \\
&= ae + bg
\end{aligned}
$$
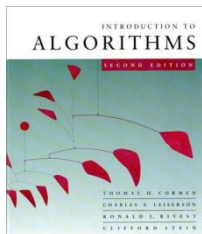
Strassen's algorithm

1. *Divide:* Partition A and B into $(n/2) \times (n/2)$ submatrices. Form P-terms to be multiplied using $+$ and $-$.

2. *Conquer:* Perform 7 multiplications of $(n/2) \times (n/2)$ submatrices recursively.

3. *Combine:* Form C using $+$ and $-$ on $(n/2) \times (n/2)$ submatrices.
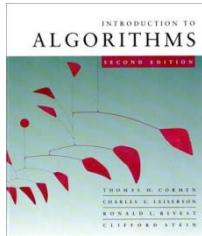
$$T(n) = 7\,T(n/2) + \Theta(n^2)$$

Analysis of Strassen

$$T(n) = 7\,T(n/2) + \Theta(n^2)$$

$$n^{\log_b a} = n^{\log_2 7} \approx n^{2.81} \implies \text{CASE 1} \implies T(n) = \Theta(n^{\log 7}).$$

The number 2.81 may not seem much smaller than 3, but because the difference is in the exponent, the impact on running time is significant. In fact, Strassen's algorithm beats the ordinary algorithm on today's machines for $n \geq 30$ or so.

Best to date (of theoretical interest only): $\Theta(n^{2.376\cdots})$.

Conclusion

- Divide and conquer is just one of several powerful techniques for algorithm design.

- Divide-and-conquer algorithms can be analyzed using recurrences and the master method (so practice this math).

- Can lead to more efficient algorithms