

## 7. Homework

Due **10/29/18** at the beginning of class

**Justify all your answers.**

### 1. Union-Find (6 points)

- (a) (3 points) Consider a disjoint-set forest data structure with path compression. Assume a sequence of MAKE-SET, UNION, and FIND-SET operations have already been performed, and the data structure contains  $n$  elements. Prove that then any consecutive sequence of FIND-SET operations takes  $O(n)$  time.
- (b) (3 points) Now consider a disjoint-set forest data structure with path compression and union-by-weight. Give a sequence of MAKE-SET, FIND-SET, and UNION operations which causes a taller tree to be linked underneath a shorter tree.

### 2. MST (9 points)

Let  $G = (V, E, w)$  be a connected undirected weighted graph.

- (a) (2 points) What are the runtimes of Prim's, Kruskal's, and Boruvka's algorithms if  $G$  is given in an adjacency matrix? How can you speed up these runtimes?
- (b) (2 points) Assume a list of all edges  $E$  is given in non-decreasing order of edge weights. In this case, what is the fastest runtime to compute an MST for  $G$ ?
- (c) (5 points) Assume for each vertex  $v \in V$ , the adjacency list  $Adj(v)$  stores adjacent vertices in non-decreasing order of edge weights. Show how to modify Boruvka's algorithm to run in  $O(|E| + |V| \log |V|)$  time in this case.

### 3. Amortized Runtime for Binary Heaps (8 points)

Consider a regular binary min-heap data structure that supports the operations INSERT and EXTRACT-MIN in  $O(\log n)$  worst-case time, where  $n$  is the number of elements in the heap. Give a potential function  $\Phi$  such that the amortized cost of INSERT is  $O(\log n)$  and the amortized cost of EXTRACT-MIN is  $O(1)$ , and show that it works.

*(Hint 1: One possible potential function is based on  $n_i \log n_i$ , where  $n_i$  is the number of elements in the heap after  $i$  operations.)*

*(Hint 2: A fact that might be useful:  $n \log \frac{n}{n-1} \leq \frac{2}{\ln 2}$ )*