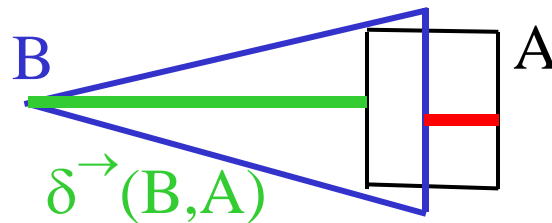
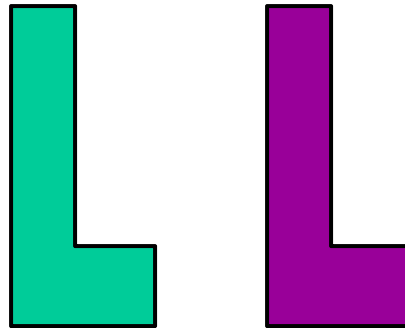
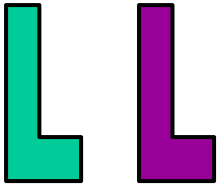

CMPS 3130/6130 Computational Geometry
Spring 2015



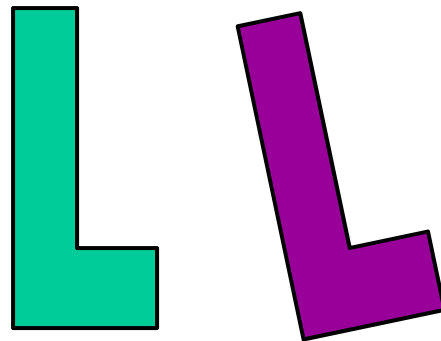
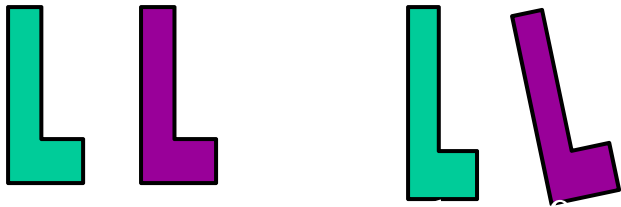
Shape Matching

Carola Wenk

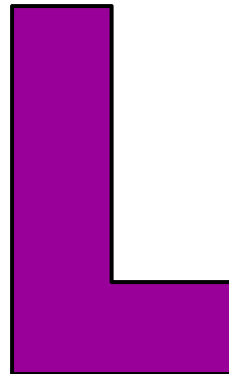
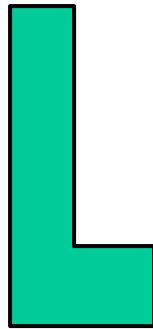
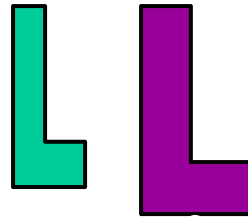
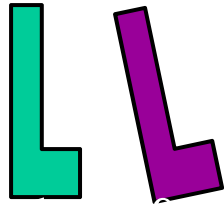
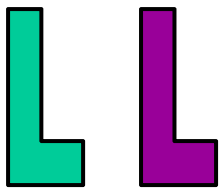
When are two shapes similar?



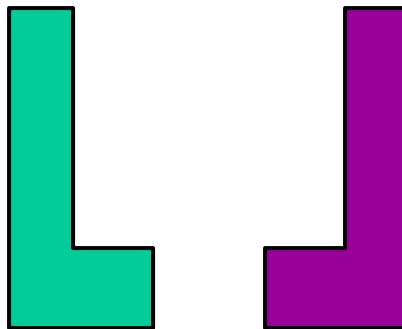
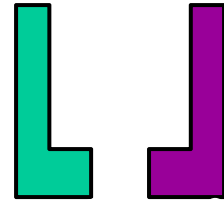
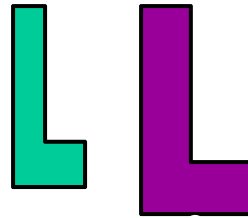
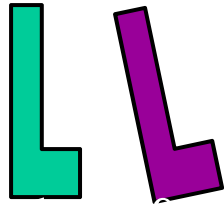
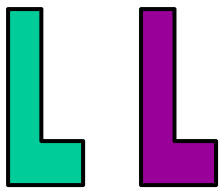
When are two shapes similar?



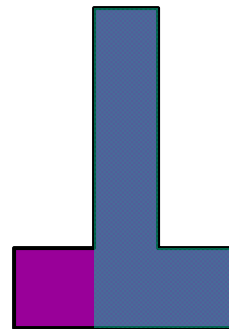
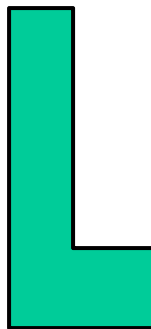
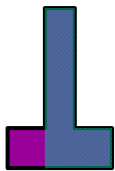
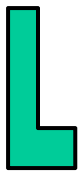
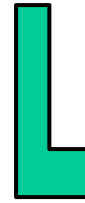
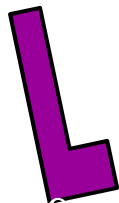
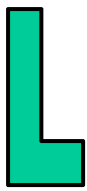
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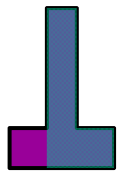
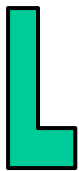
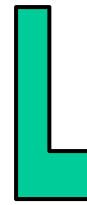
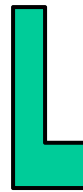
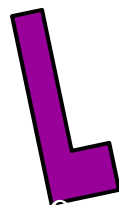
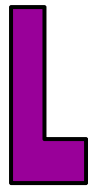
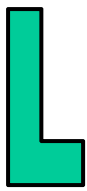
When are two shapes similar?



When are two shapes similar?



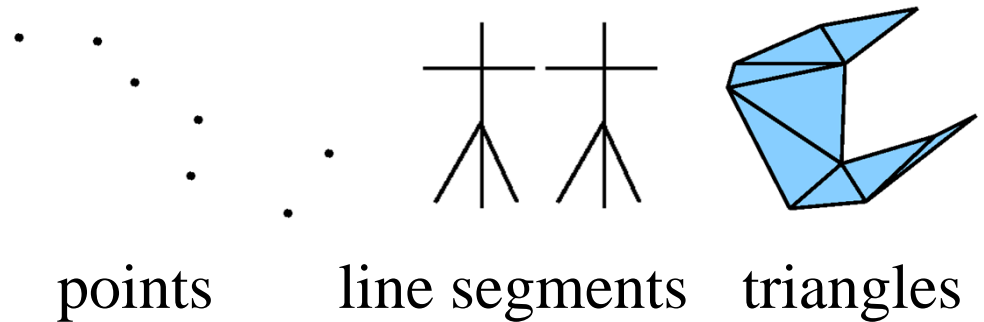
When are two shapes similar?



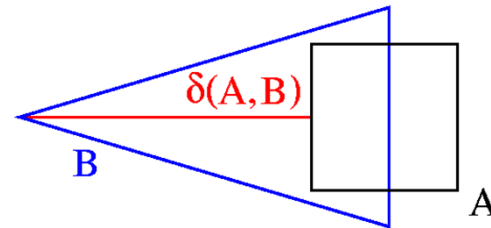
Geometric Shape Matching

Given:

Two **geometric shapes**, each composed of a number of basic objects such as



An application dependant **distance measure δ** , e.g., the Hausdorff distance



A set of **transformations T** , e.g., translations, rigid motions, or none

Matching Task:

Compute $\min_{T \in T} \delta(T(A), B)$

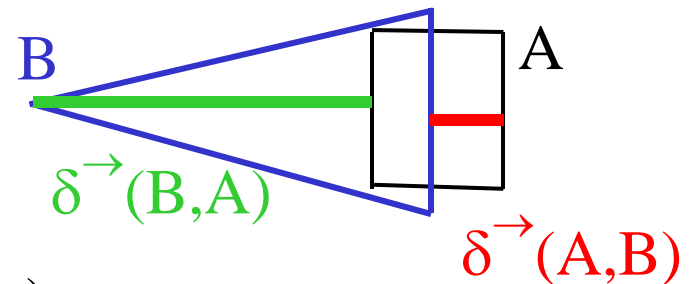
Hausdorff distance

- Directed Hausdorff distance

$$\delta^{\rightarrow}(A,B) = \max_{a \in A} \min_{b \in B} \| a-b \|$$

- Undirected Hausdorff-distance

$$\delta(A,B) = \max (\delta^{\rightarrow}(A,B) , \delta^{\rightarrow}(B,A))$$



- If A, B are discrete point sets, $|A|=m, |B|=n$
 - In d dimensions: Compute in $O(mn)$ time brute-force
 - In 2 dimensions: Compute in $O((m+n) \log (m+n))$ time
 - Compute $\delta^{\rightarrow}(A,B)$ by computing VD(B) and performing nearest neighbor search for every point in A

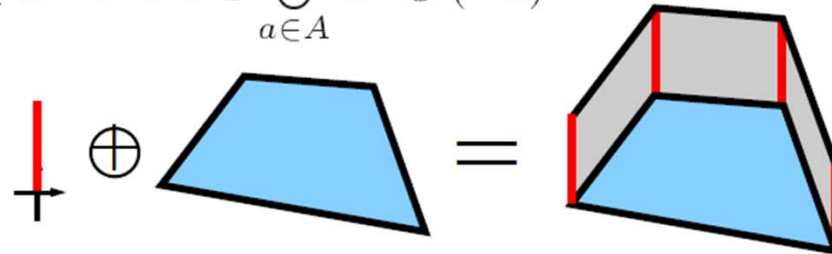
Hausdorff distance under translations

Decision problem, transformation space approach:

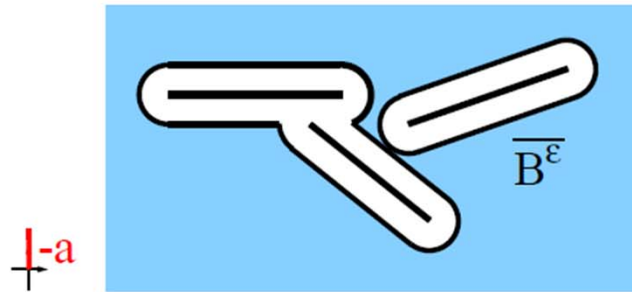
$$\vec{\delta}_H(A + t, B) \leq \epsilon \iff A + t \subseteq B^\epsilon \iff t \in \bigcup_{a \in A} \overline{B^\epsilon} \oplus (-a)$$

- Minkowski sum

$$a \oplus b = \bigcup_{x \in a} \bigcup_{y \in b} x + y$$



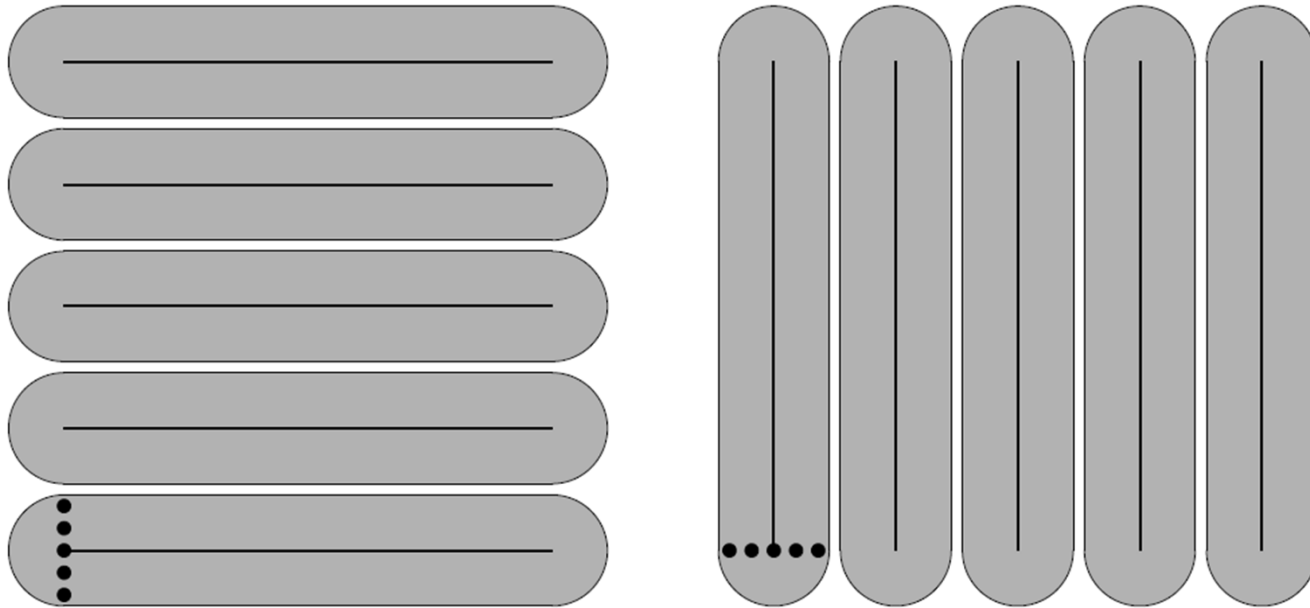
- Arrangement construction



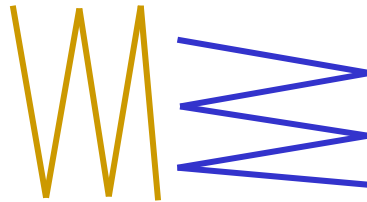
- Segments, \mathbb{R}^2 : $O((n \cdot n)^2)$ arrangement complexity (AST94)
- Simplices, \mathbb{R}^d , L_∞ : $O((\text{Complexity}(B^\epsilon))^d n^d) \leq O(n^{d^2+d})$
- Triangles, \mathbb{R}^3 , L_2 : $O(n^{6+\delta})$ arrangement complexity

Hausdorff distance

- $\Omega(m^2n^2)$ lower bound construction for directed Hausdorff distance of line segments under translations

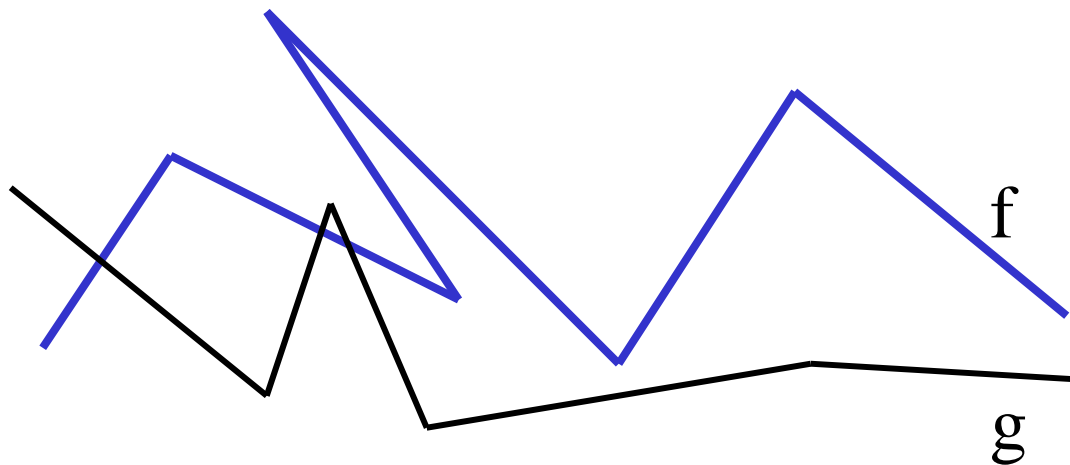


Comparing Curves and Trajectories



Polygonal Curves

- Let $f, g: [0, 1] \rightarrow \mathbb{R}^d$ be two **polygonal curves** (i.e., piecewise linear curves)



- What are good distance measures for curves?
 - Hausdorff distance?
 - Fréchet distance?

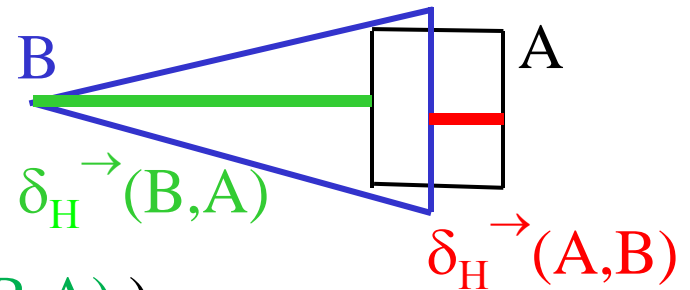
When Are Two Curves „Similar“?

- Directed Hausdorff distance

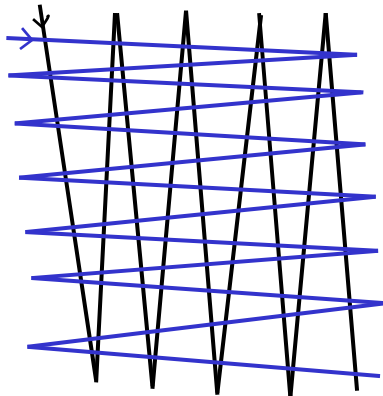
$$\delta_H^{\rightarrow}(A,B) = \max_{a \in A} \min_{b \in B} \| a-b \|$$

- Undirected Hausdorff-distance

$$\delta_H(A,B) = \max(\delta_H^{\rightarrow}(A,B), \delta_H^{\rightarrow}(B,A))$$



But:

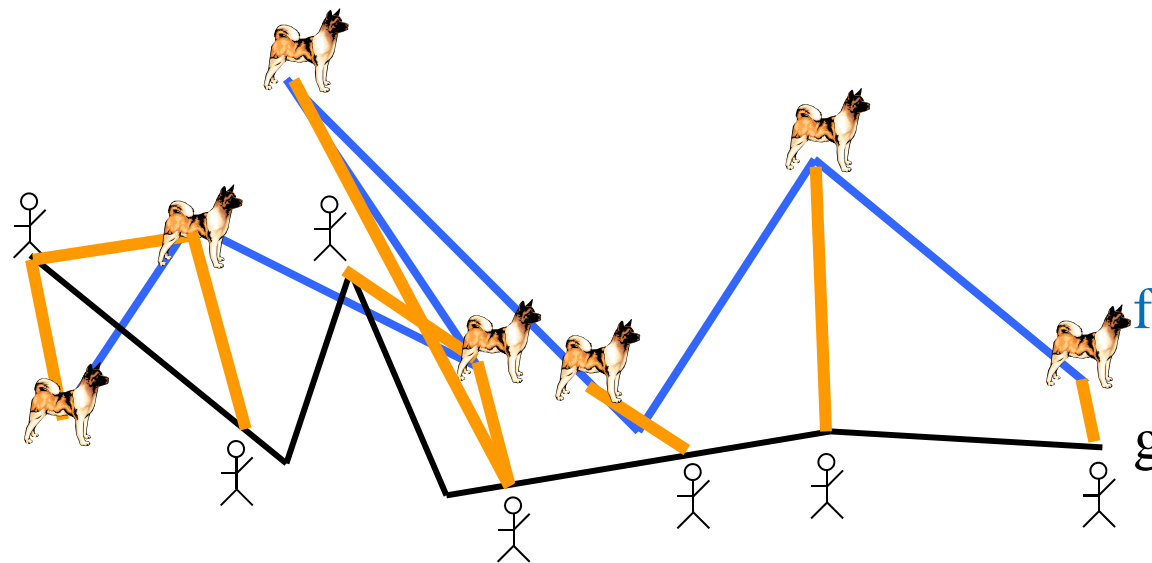


- Small Hausdorff distance
- When considered as curves the distance should be large
- The Fréchet distance takes the continuity of the curves into account

Fréchet Distance for Curves

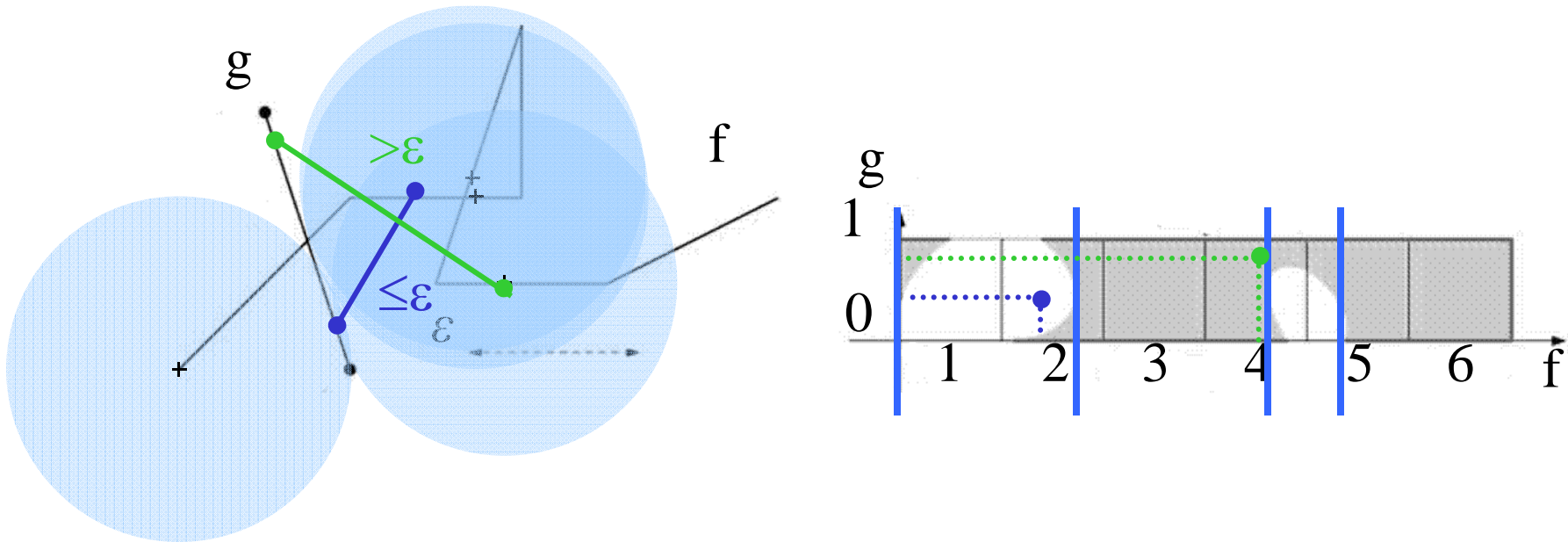
$$\delta_F(f,g) = \inf_{\alpha, \beta: [0,1] \rightarrow [0,1]} \max_{t \in [0,1]} \|f(\alpha(t)) - g(\beta(t))\|$$

where α and β range over continuous monotone increasing reparameterizations only.



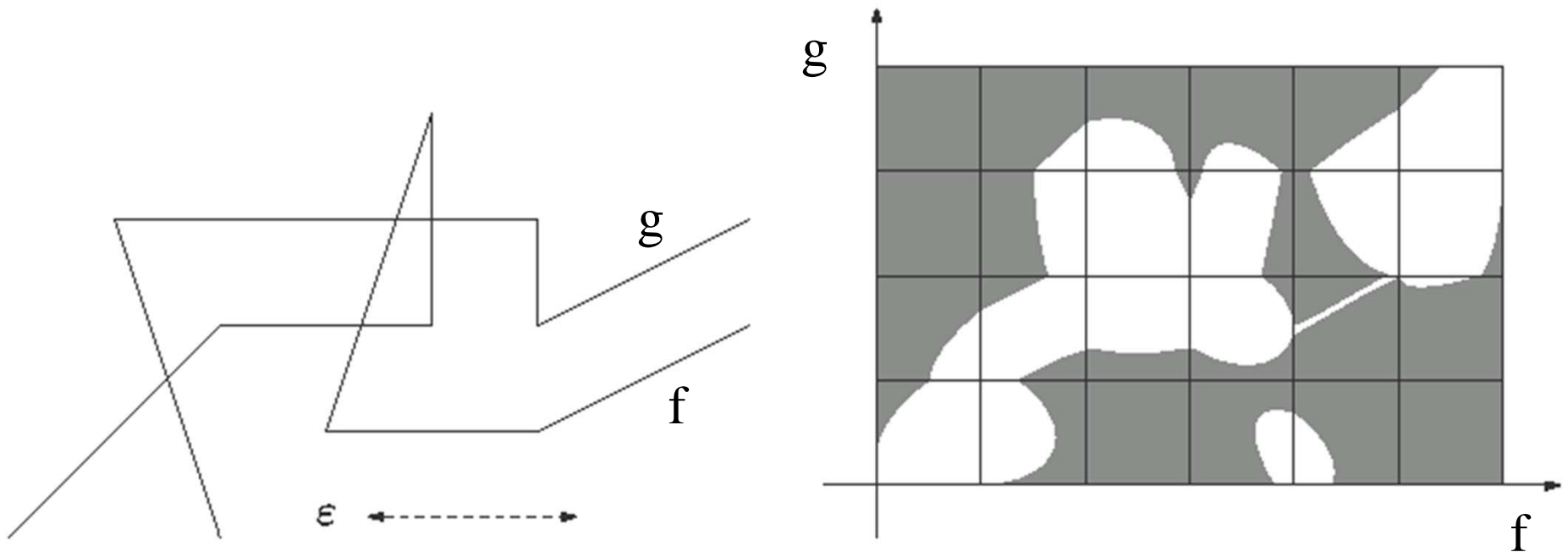
- Man and dog walk on one curve each
- They hold each other at a **leash**
- They are only allowed to go forward
- δ_F is the minimal possible leash length

Free Space Diagram



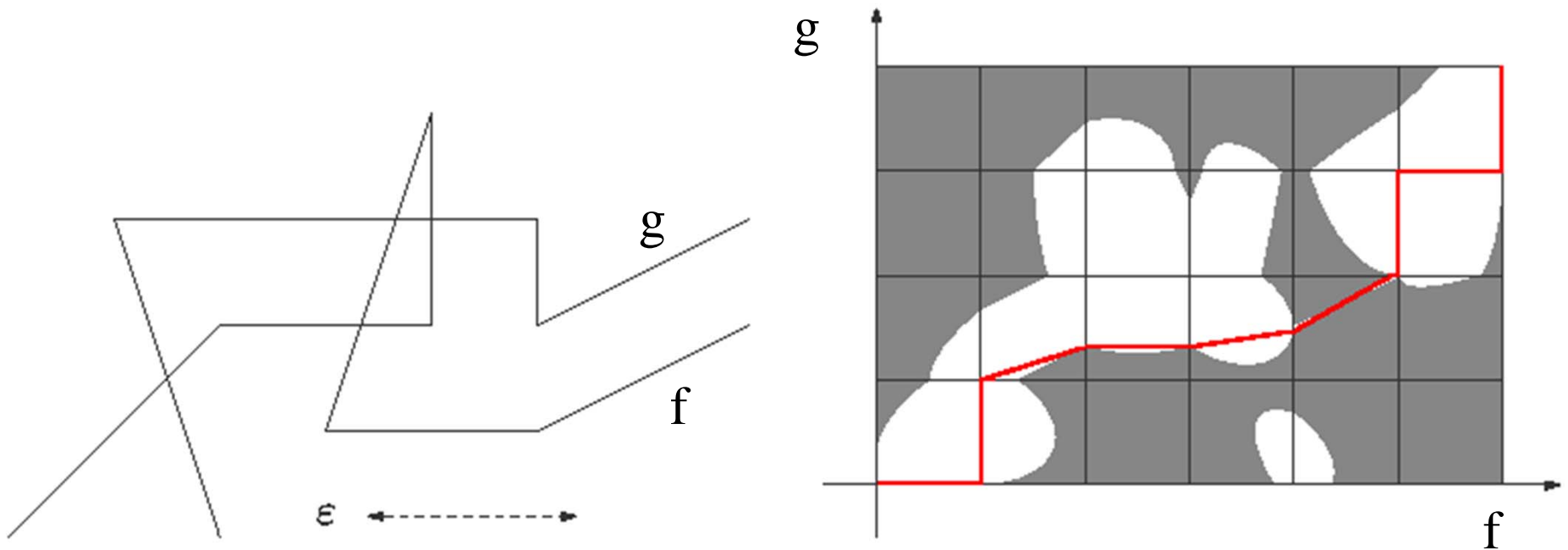
- Let $\varepsilon > 0$ fixed (eventually solve decision problem)
- $F_\varepsilon(f, g) = \{ (s, t) \in [0, 1]^2 \mid \|f(s) - g(t)\| \leq \varepsilon \}$ *white points*
free space of f and g
- The free space in one cell is an ellipse.

Free Space Diagram



- Let $\epsilon > 0$ fixed (eventually solve decision problem)
- $F_\epsilon(f, g) = \{ (s, t) \in [0, 1]^2 \mid \| f(s) - g(t) \| \leq \epsilon \}$ *white points*
free space of f and g
- The free space in one cell is an ellipse.

Free Space Diagram



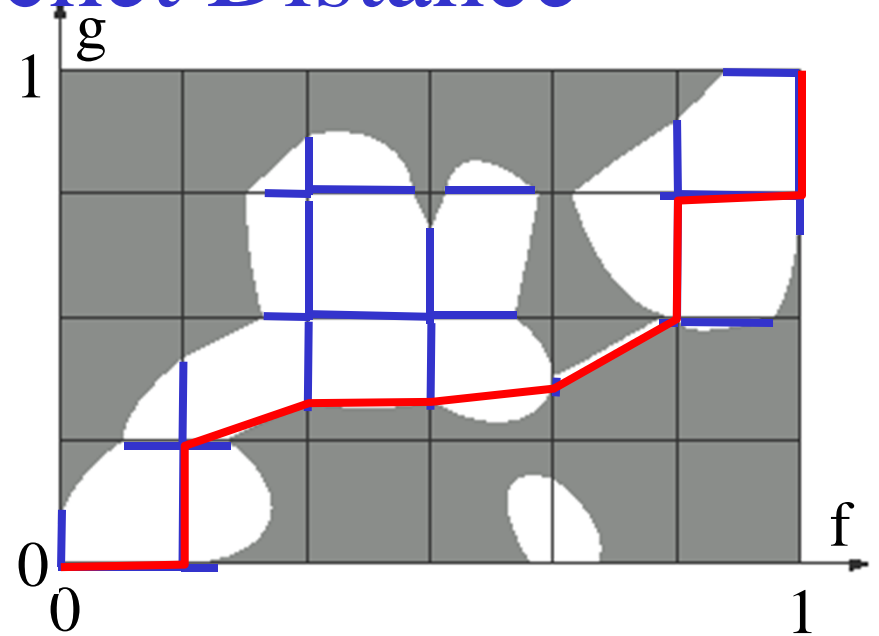
- Monotone path encodes reparametrizations of f and g
- $\delta_F(f,g) \leq \varepsilon$ iff there is a monotone path in the free space from $(0,0)$ to $(1,1)$

Compute the Fréchet Distance

- **Solve the decision problem**

$\delta_F(f,g) \leq \varepsilon$ in $O(mn)$ time:

- Find monotone path using DP:
- On each cell boundary compute the interval of all points that are reachable by a monotone path from (0,0)
- Compute a **monotone path** by backtracking



- **Solve the optimization problem**

- In practice in $O(mn \log b)$ time with binary search and b-bit precision
- In $O(mn \log mn)$ time [AG95] using parametric search (using Cole's sorting trick)
- In $O(mn \log^2 mn)$ expected time [CW09] with randomized red/blue intersections

[AG95] H. Alt, M. Godau, Computing the Fréchet distance between two polygonal curves, *IJCGA* 5: 75-91, 1995.

[CW10] A.F. Cook IV, C. Wenk, Geodesic Fréchet Distance Inside a Simple Polygon, *ACM TALG* 7(1), 19 pages, 2010.

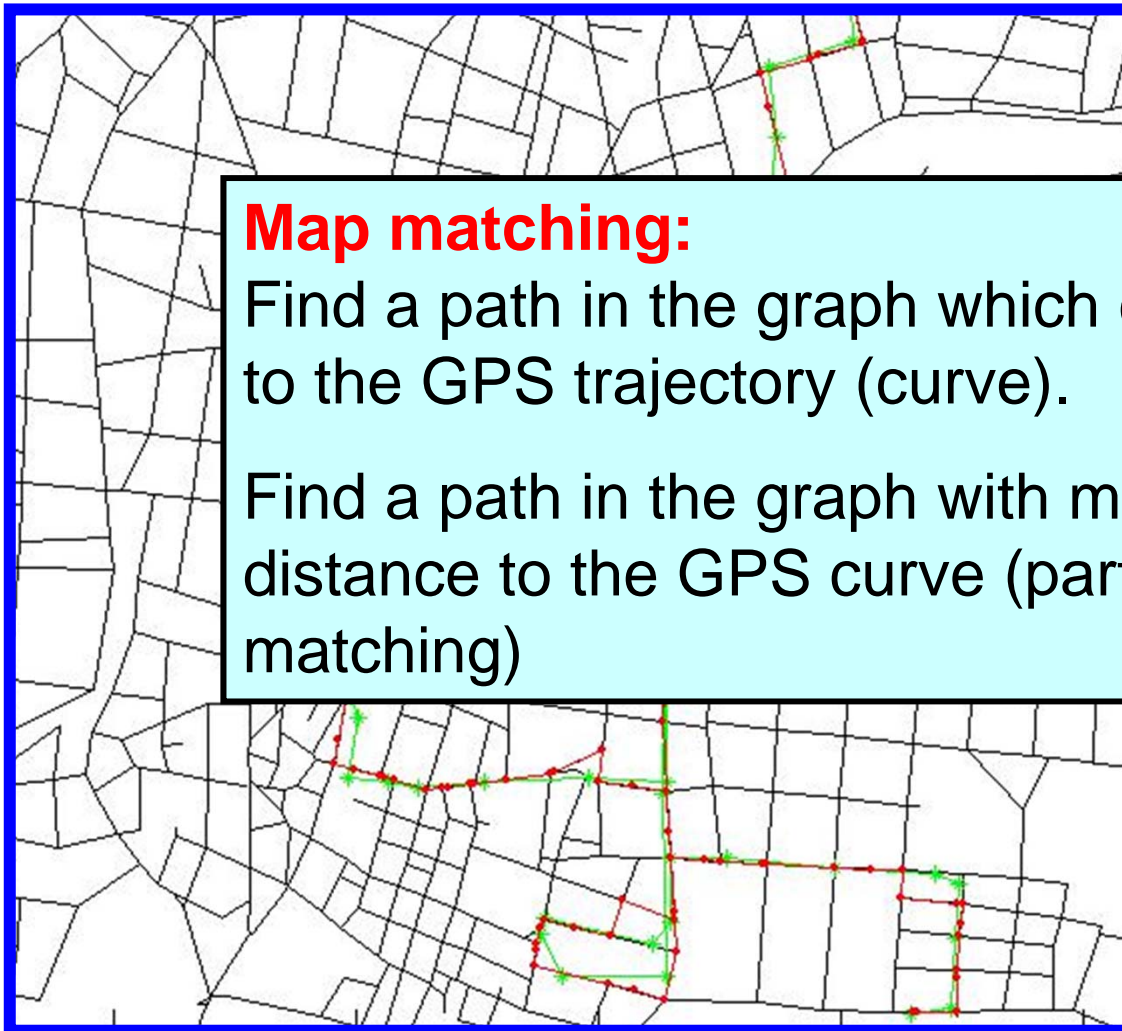


GPS Trajectories for Dynamic Routing

- Navigation systems answer shortest path queries based on travel times on road segments
- Usually based on static travel-time weights derived from speed limits
- **Goal:** Develop a navigation system which answers routing queries **based on the current traffic situation**
- **Current traffic situation:**
 - Database of current travel times per road segment
 - Use GPS trajectory data from vehicle fleets



GPS Trajectories from Vehicles



1) **Measurement error:**
GPS points

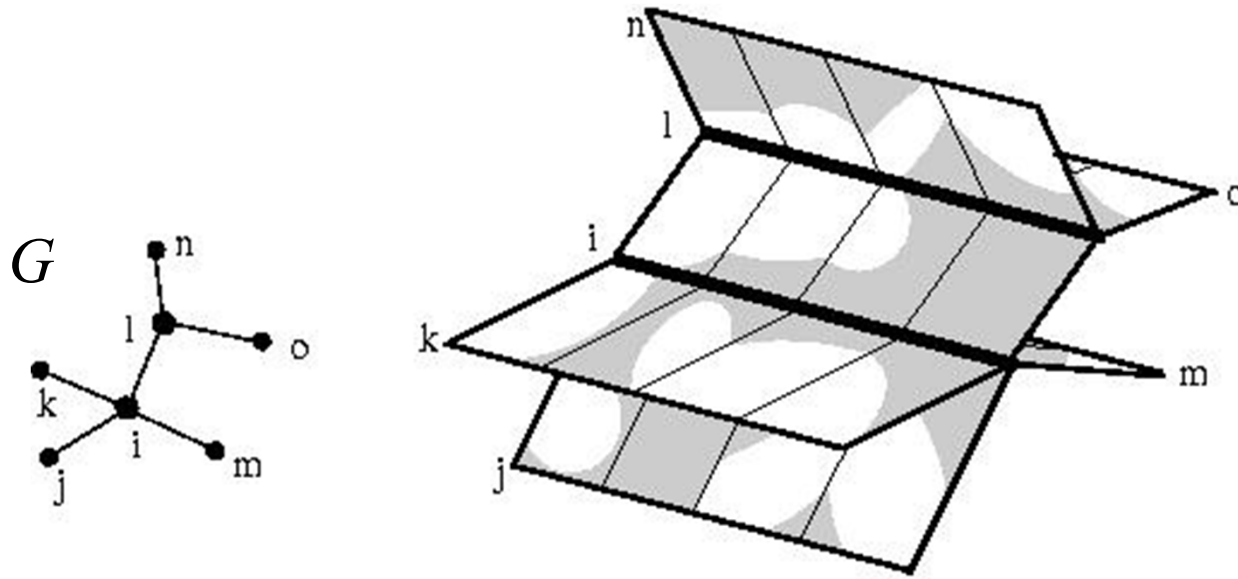
do not lie
d map

error:
trajectory
product and
sampled

every 50s

⇒ The GPS trajectory
does not lie on the
road map

Free Space Surface



- Glue the free space diagrams $FD_{i,j}$ together according to adjacency information in G

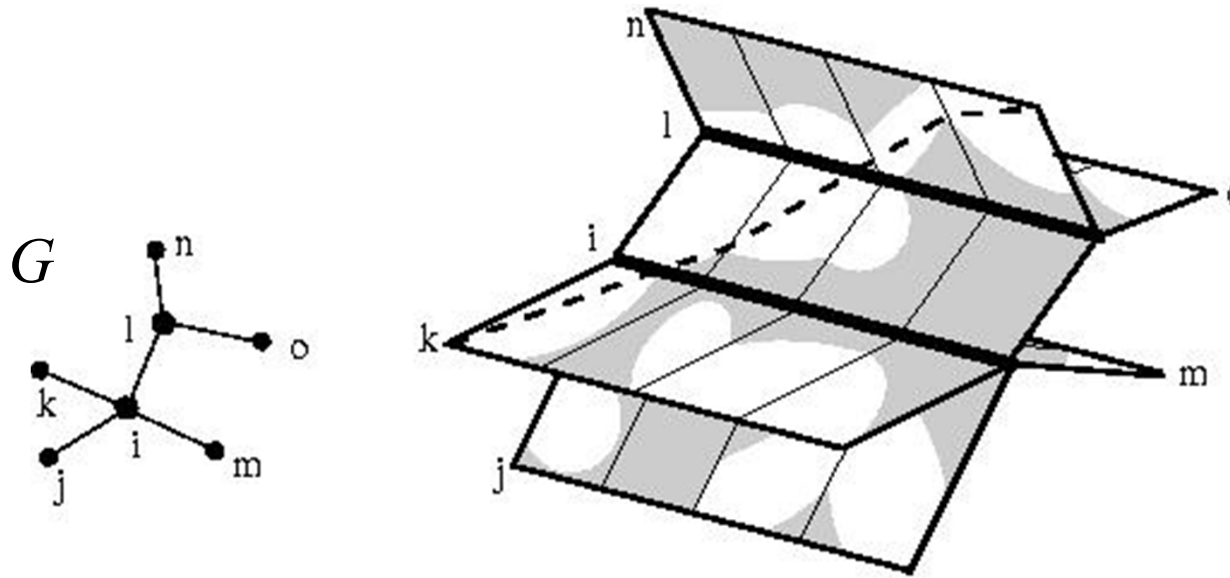
➔ **Free space surface of f and G**

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Free Space Surface



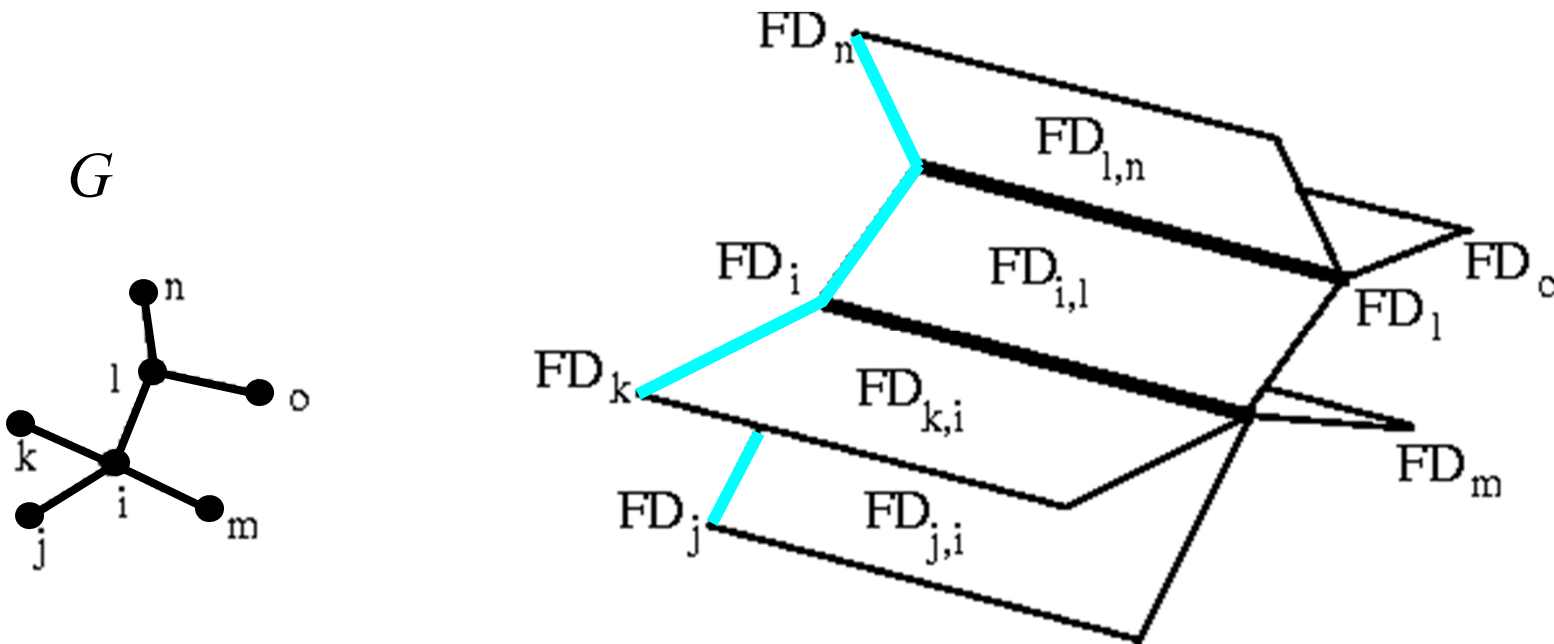
- **Task:** Find a **monotone path** in the free space surface that
 - starts in a lower left corner
 - and ends in an upper right corner

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Sweep



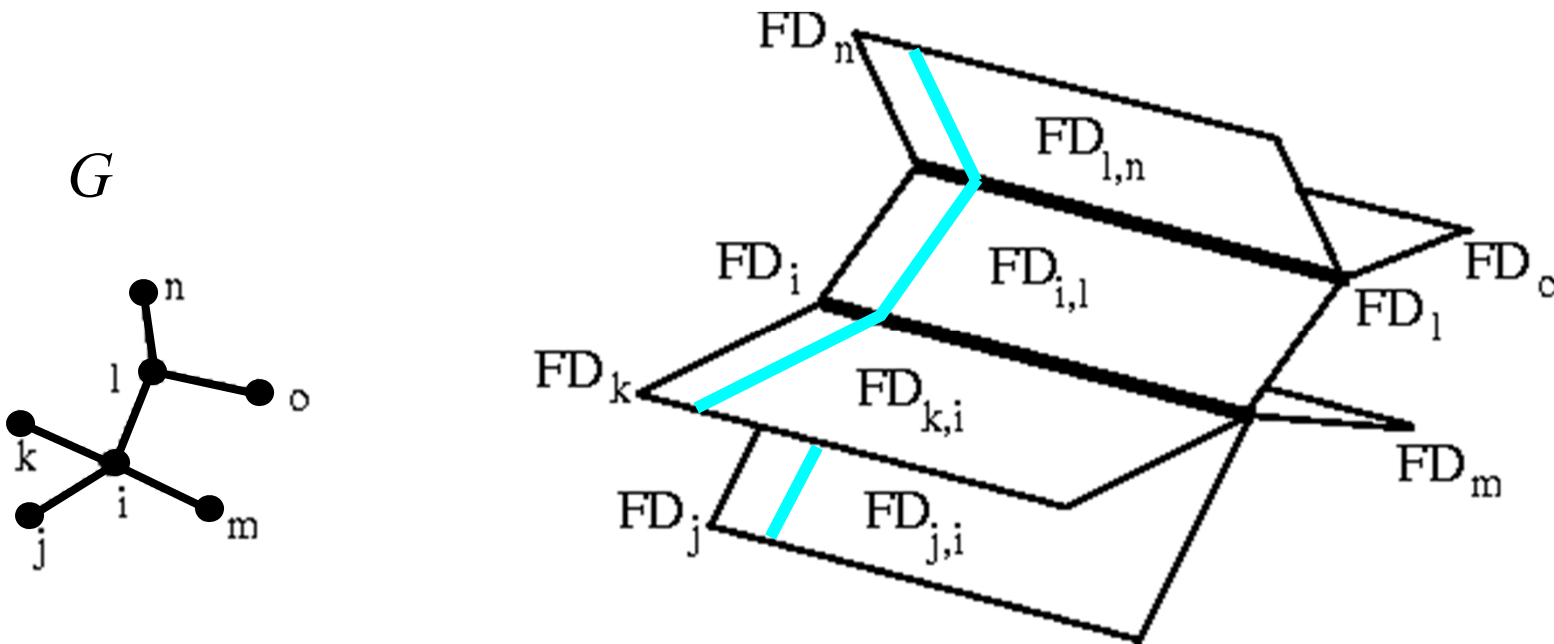
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

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Sweep



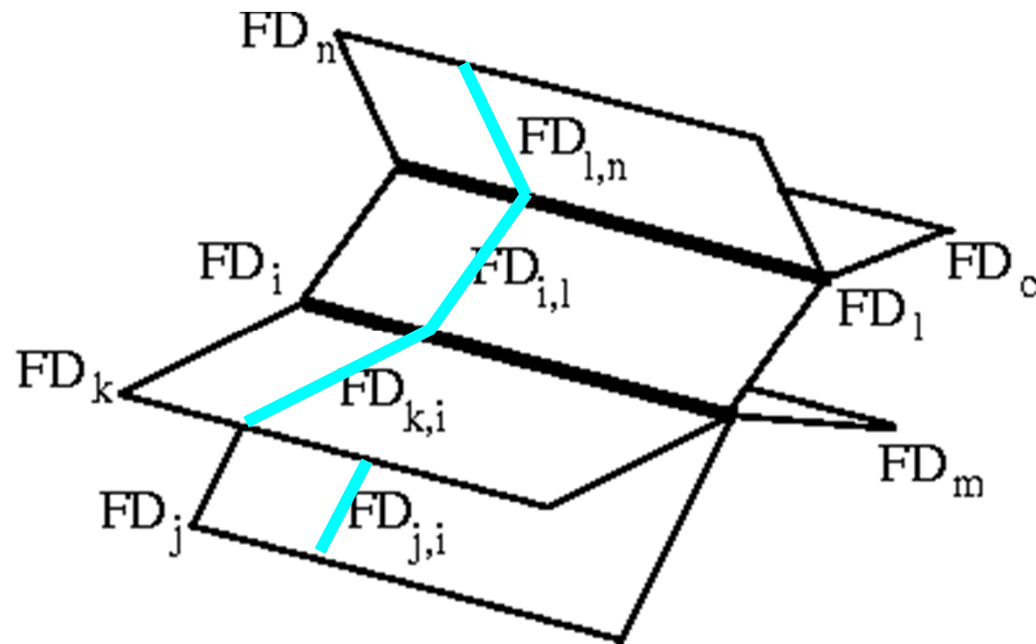
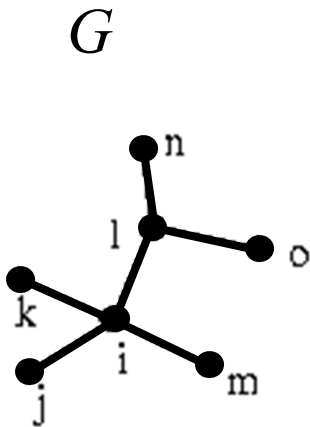
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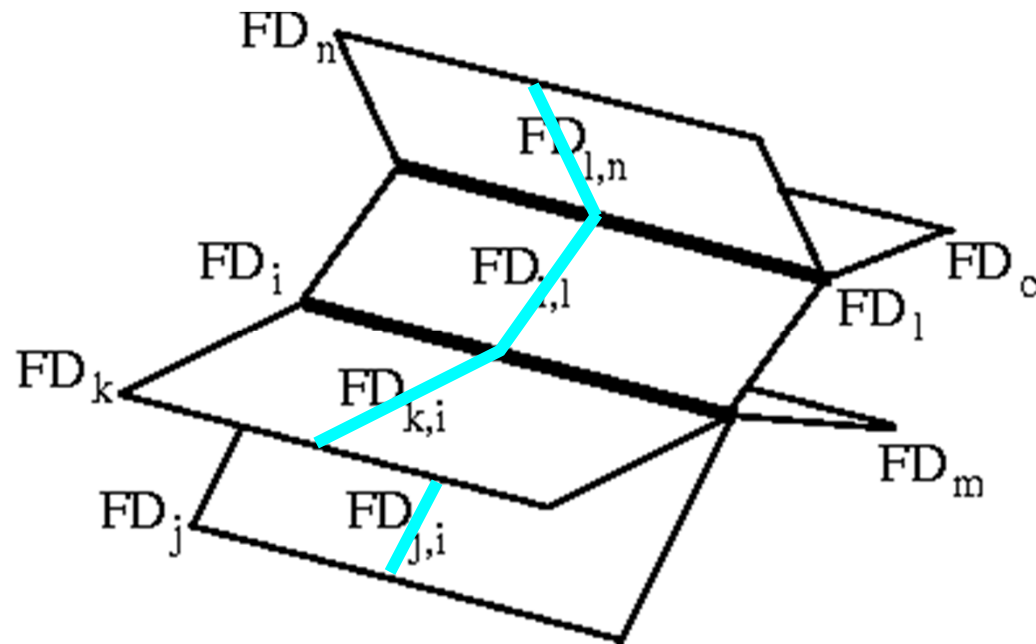
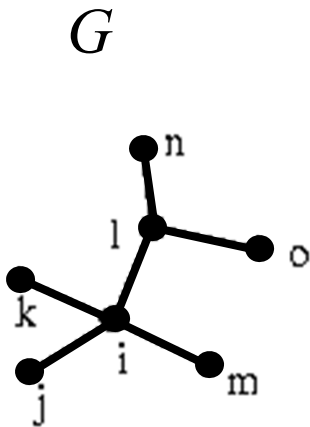
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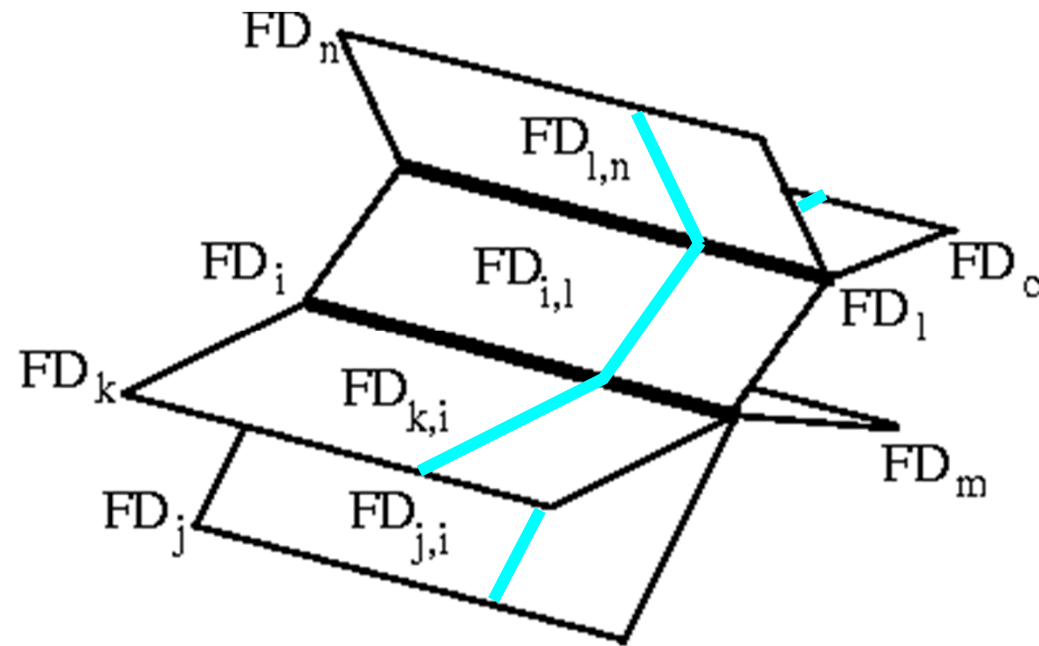
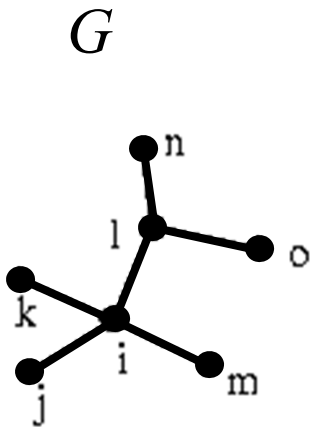
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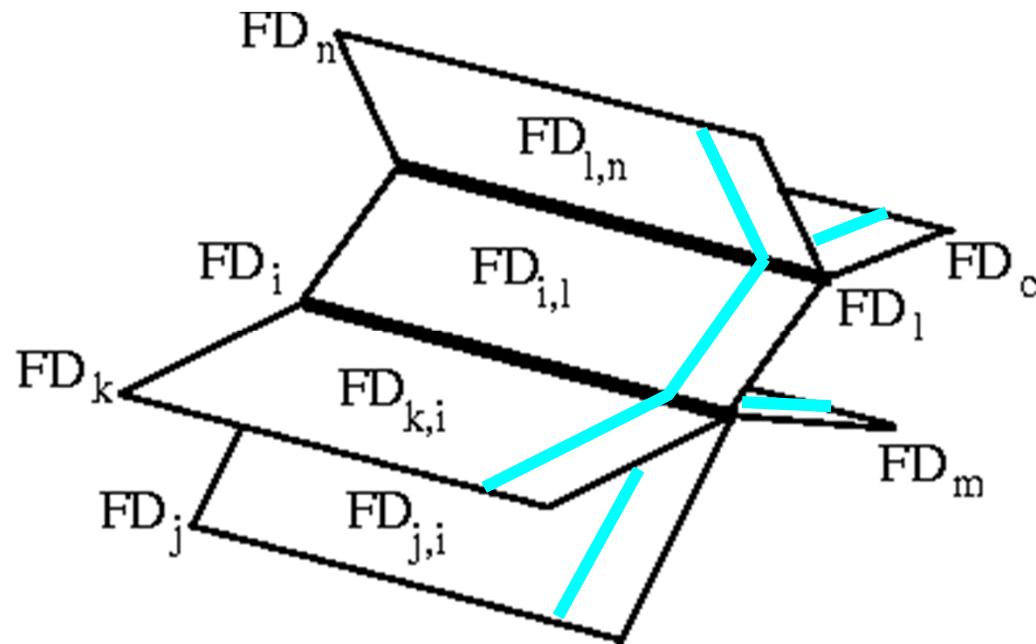
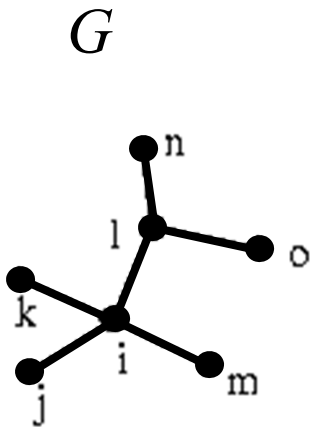
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Sweep



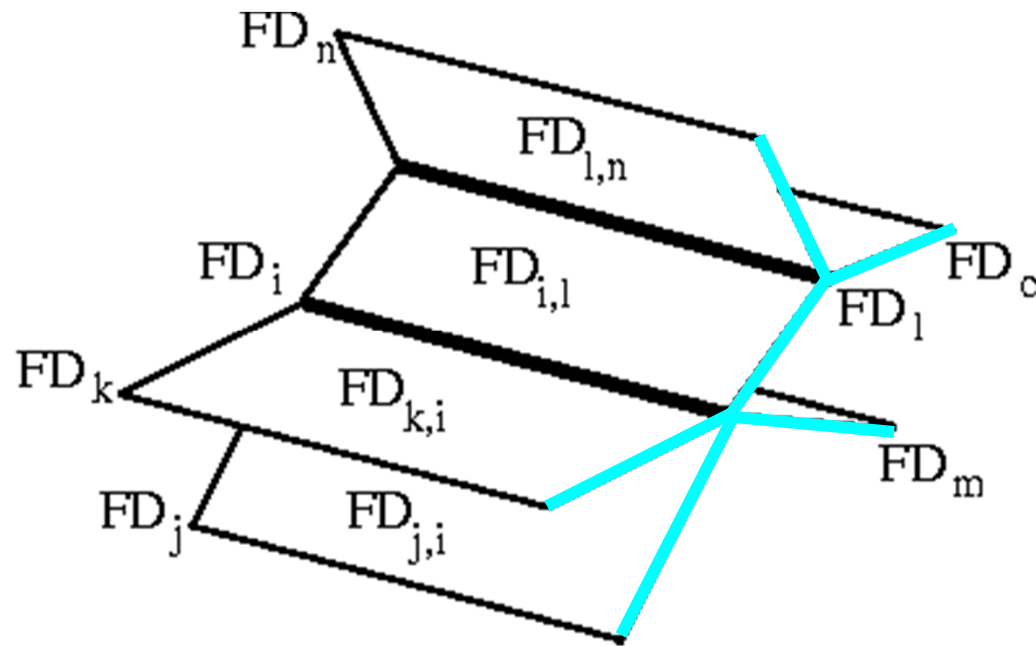
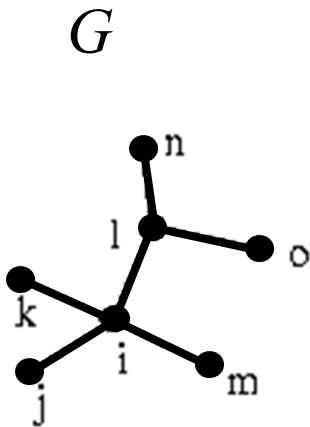
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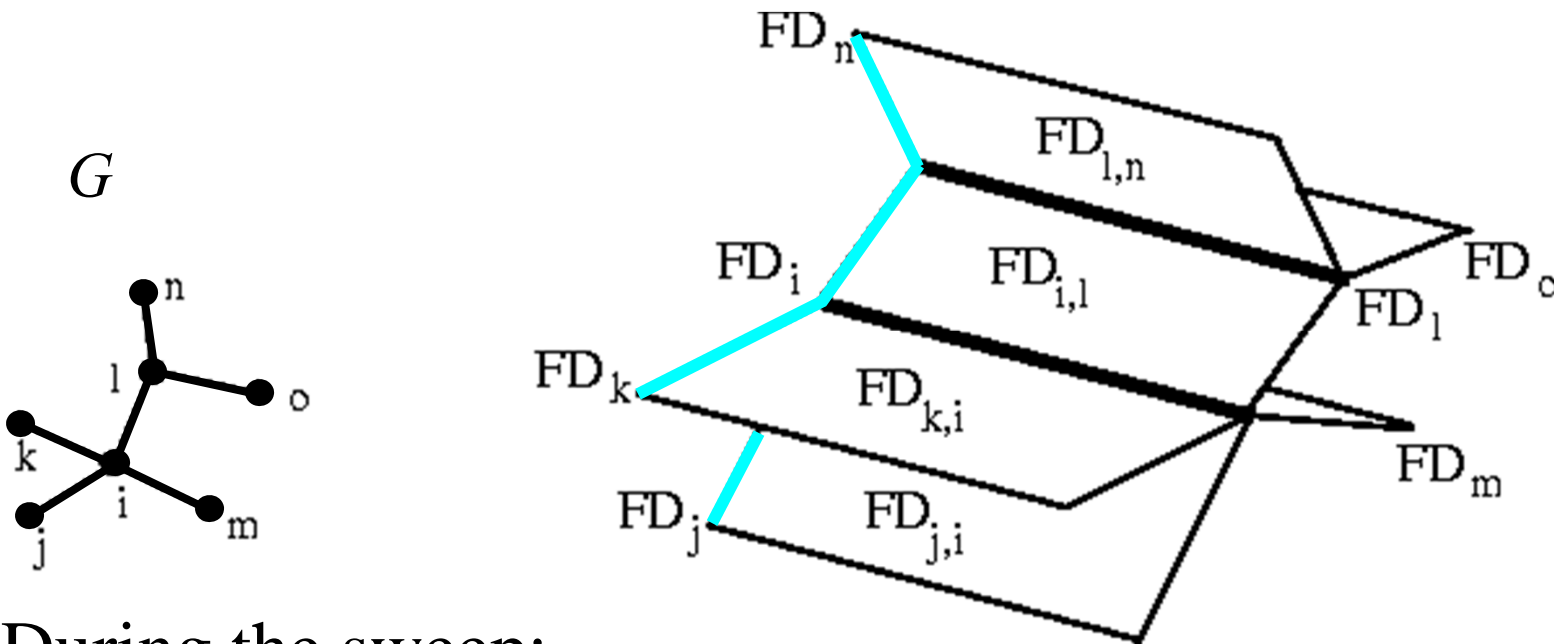
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[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Compute Reachable Points



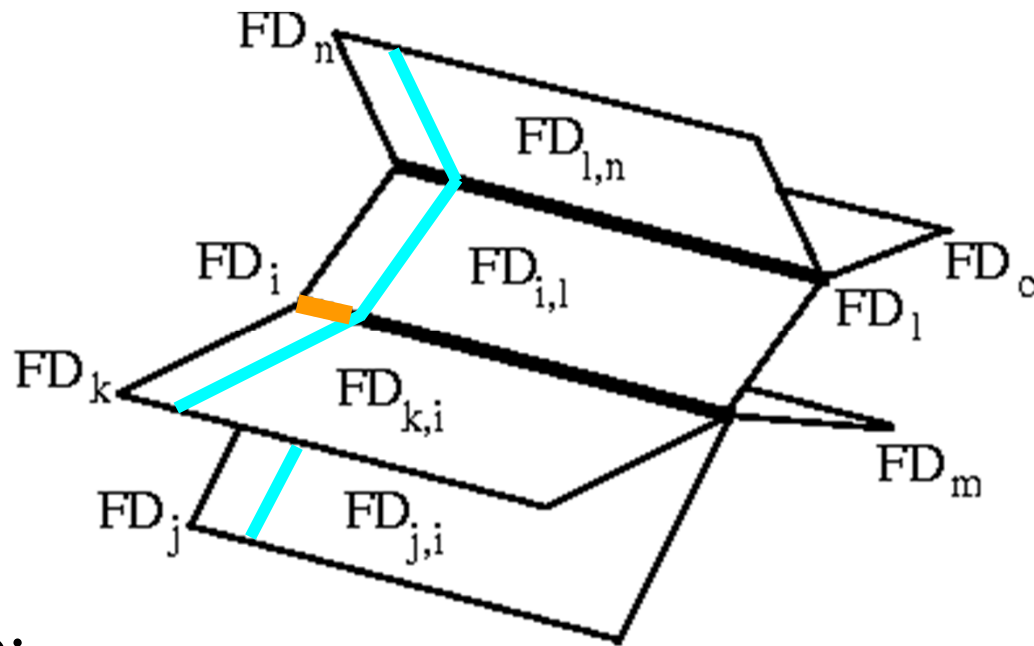
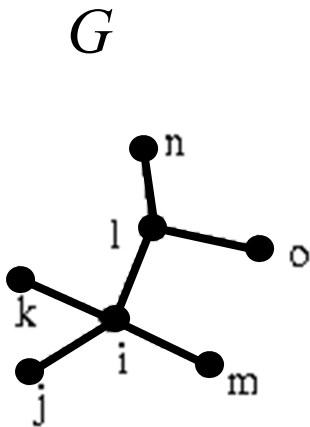
- During the sweep:
Compute **points** on the free space surface, to the left of the sweep line, which are reachable by a monotone path from a lower left corner.

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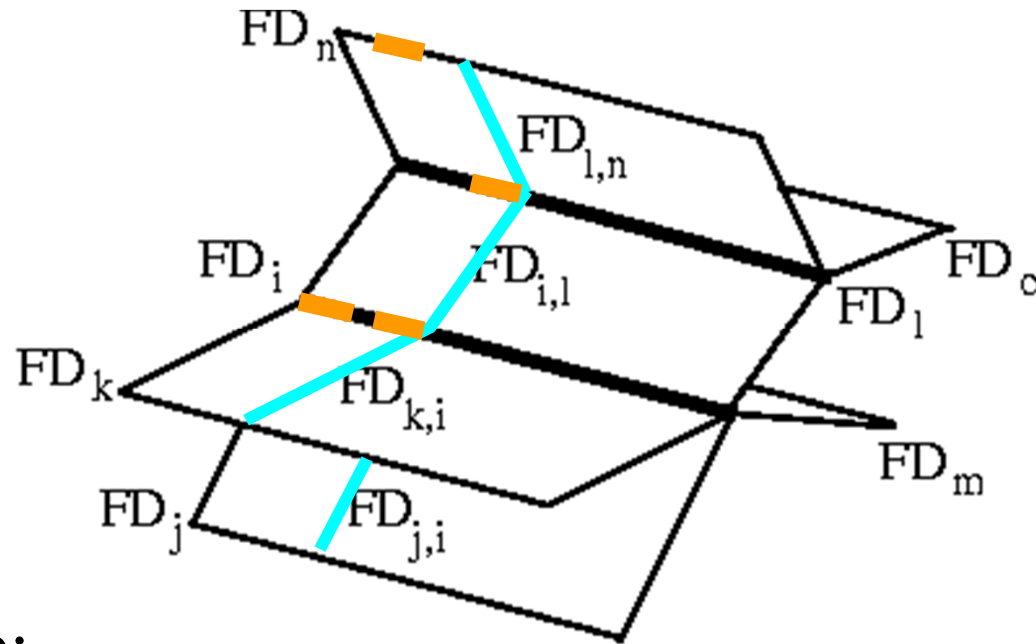
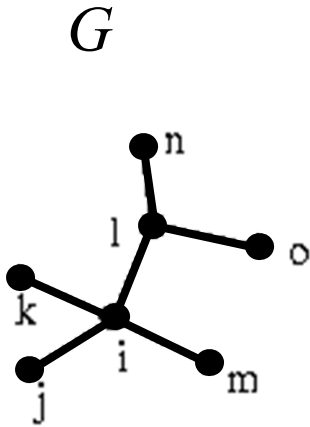
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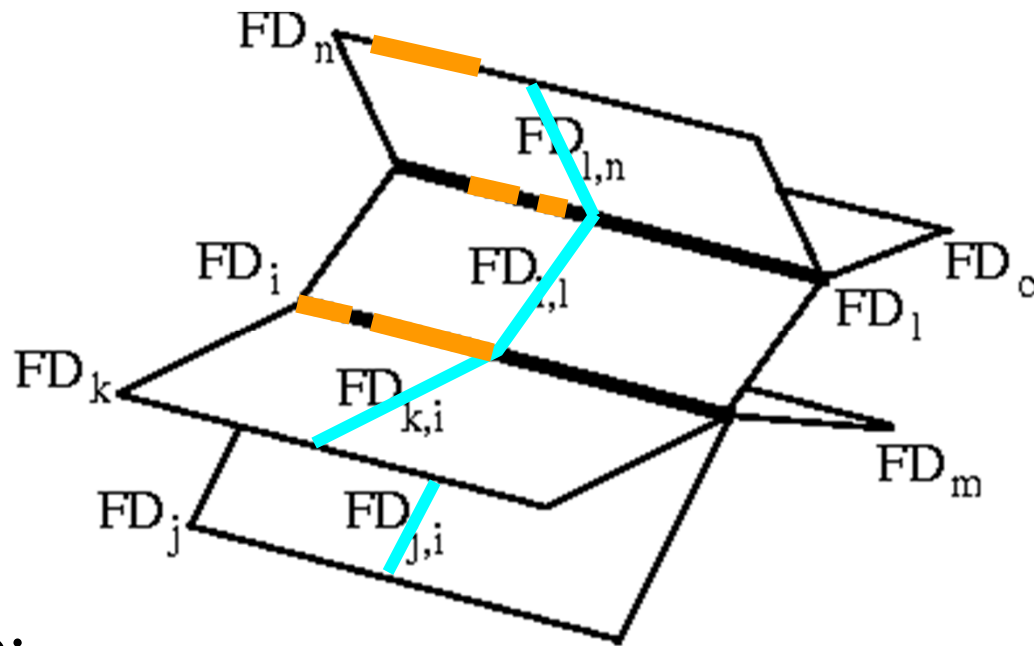
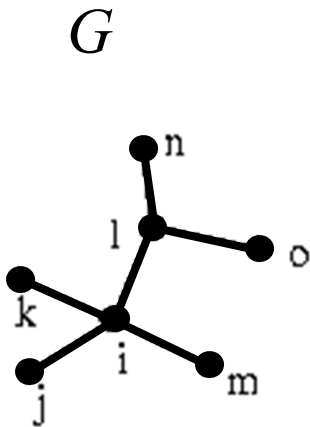
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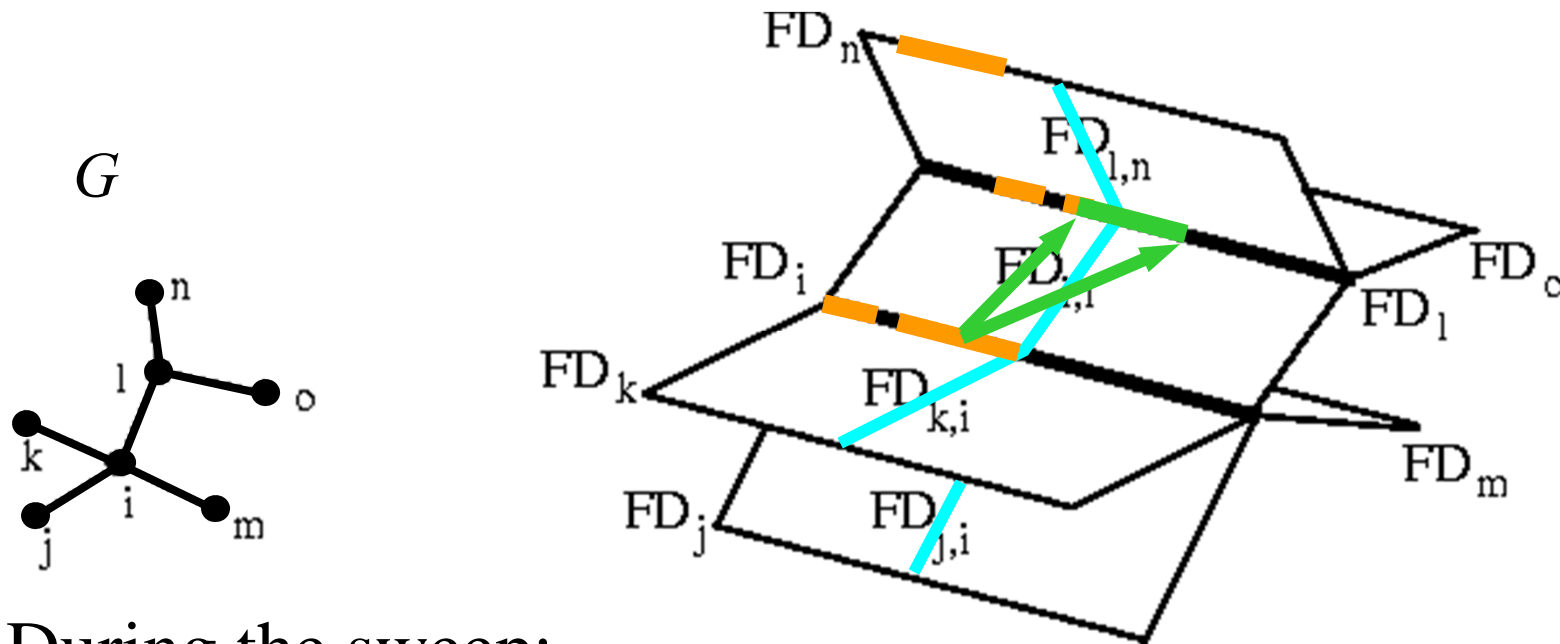
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Update Reachable Points



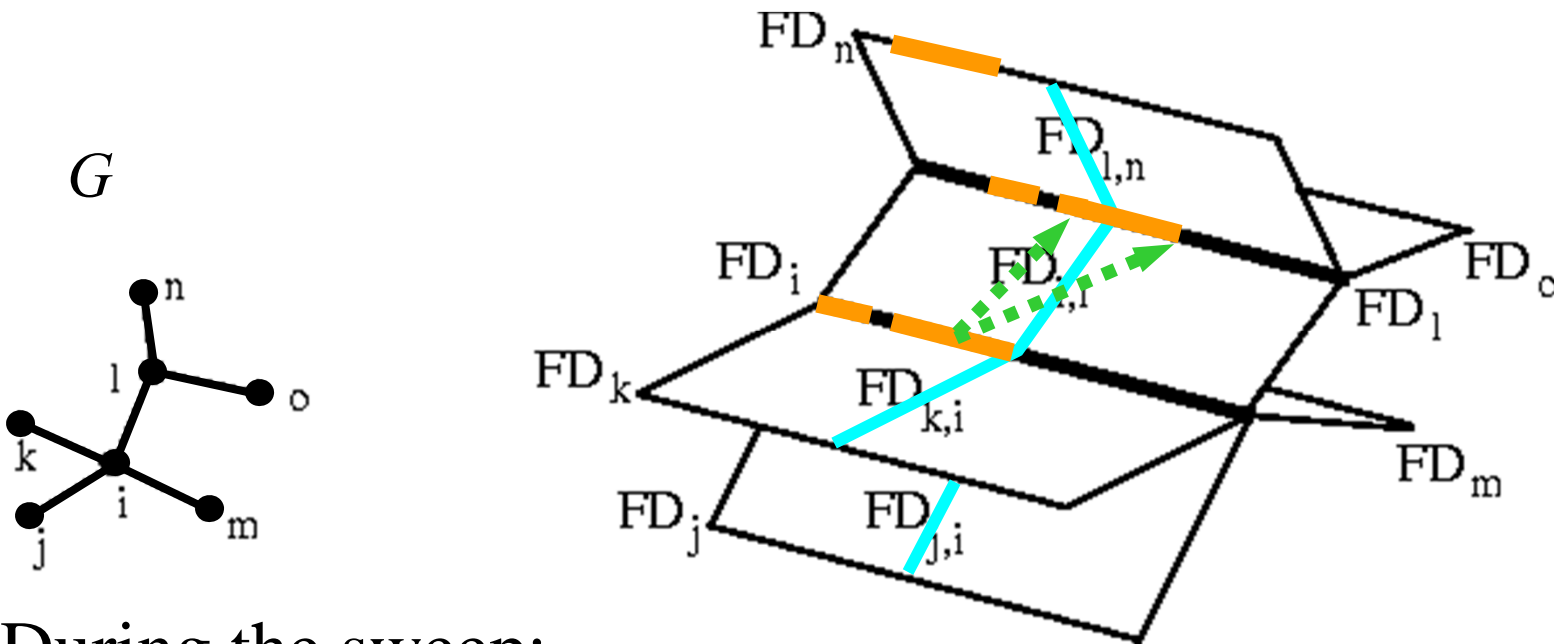
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 - Use a data structure which supports reachability queries in the free space surface

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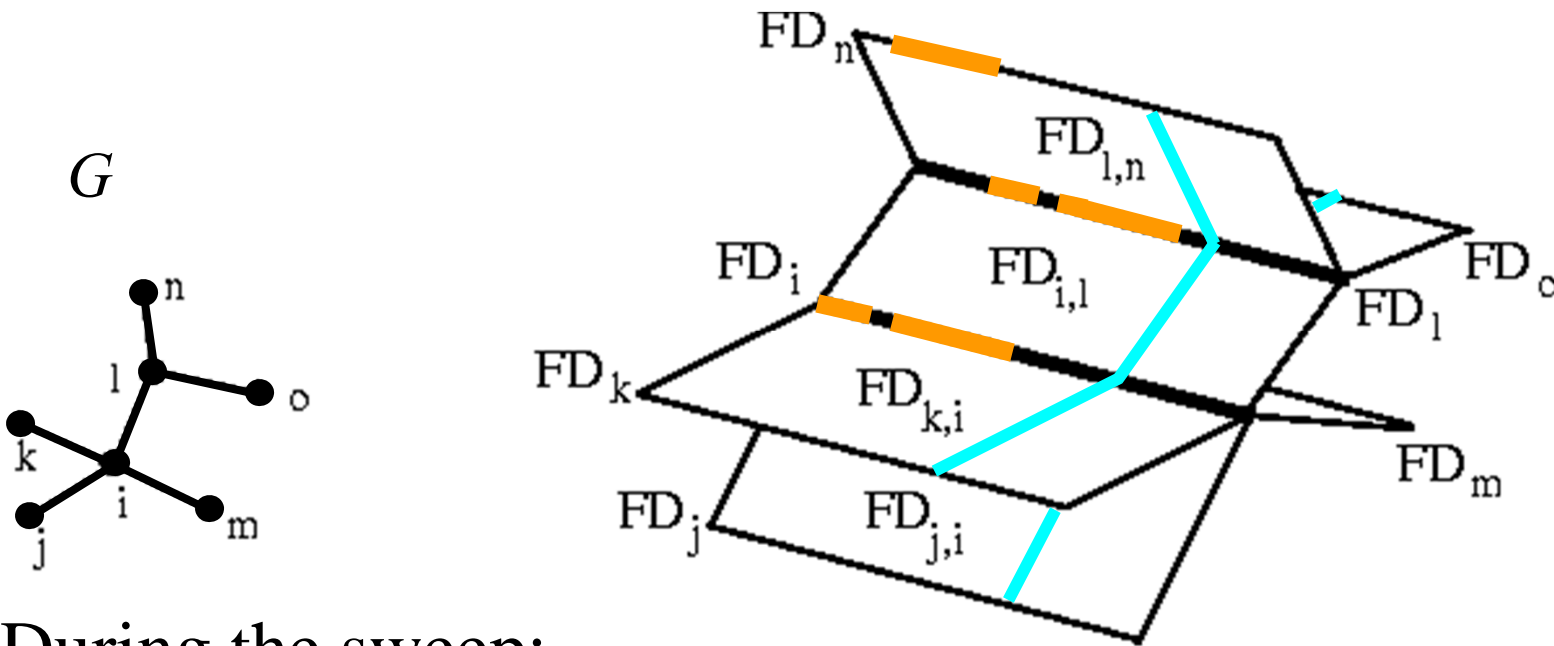
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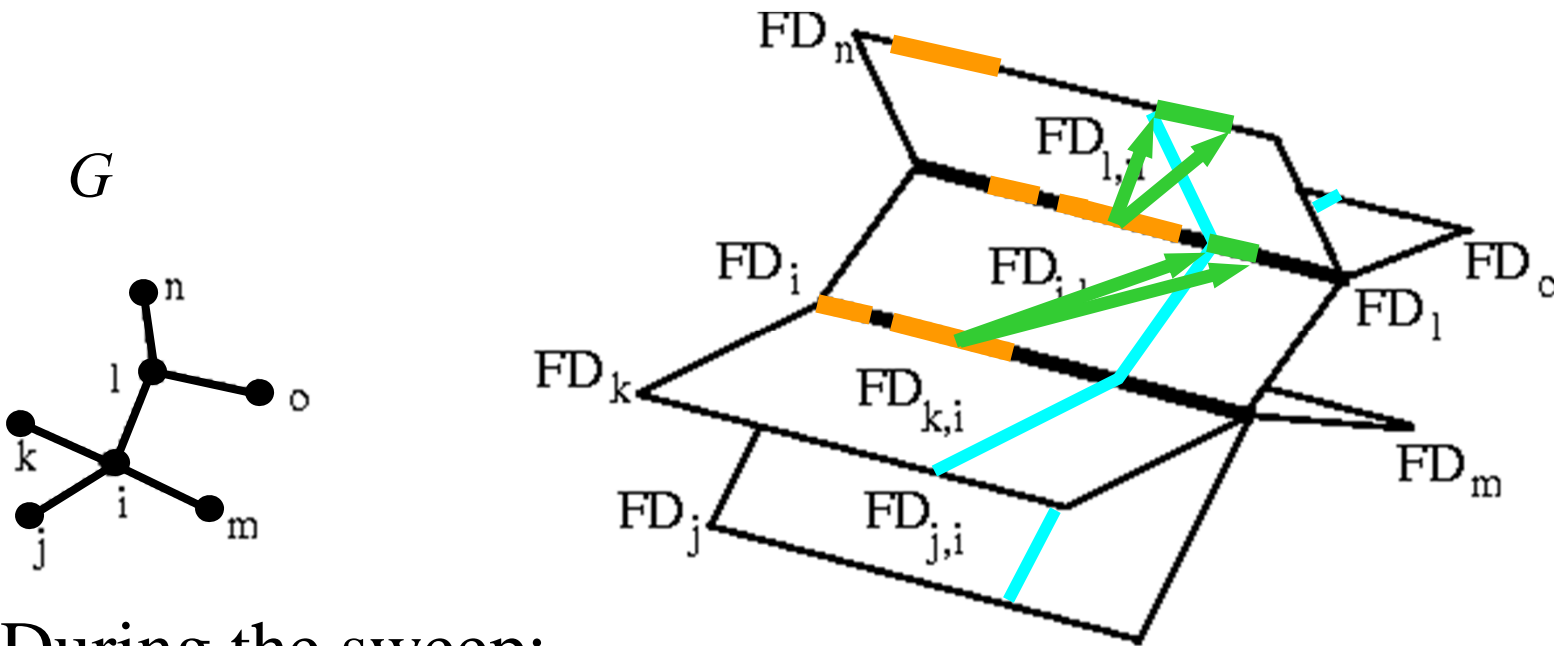
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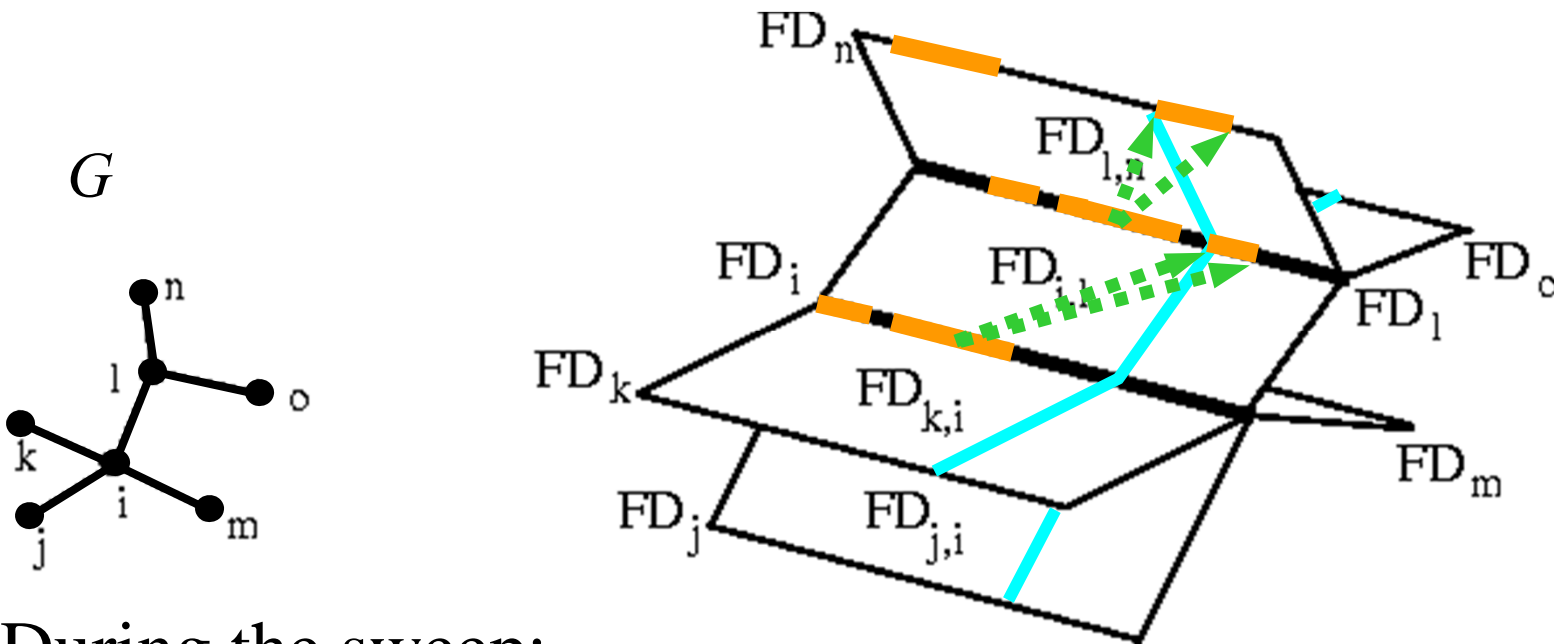
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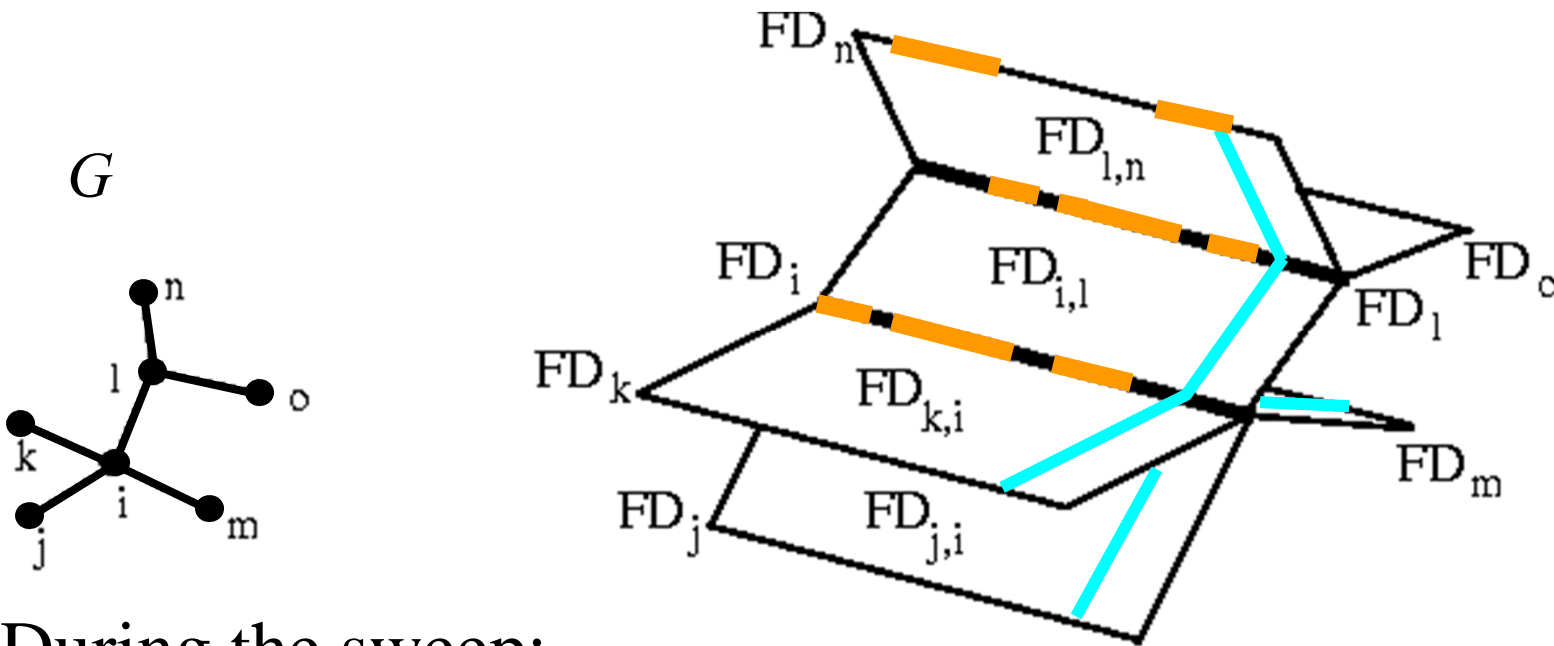
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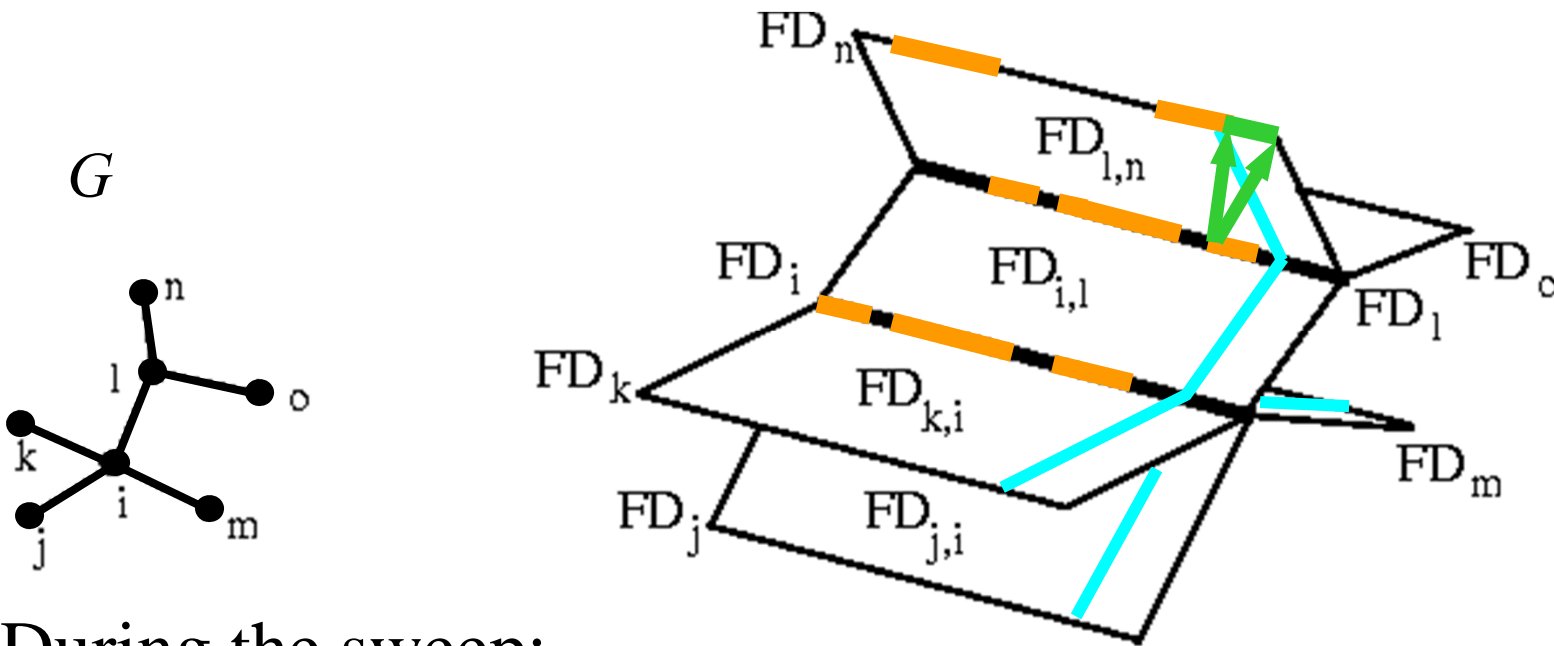
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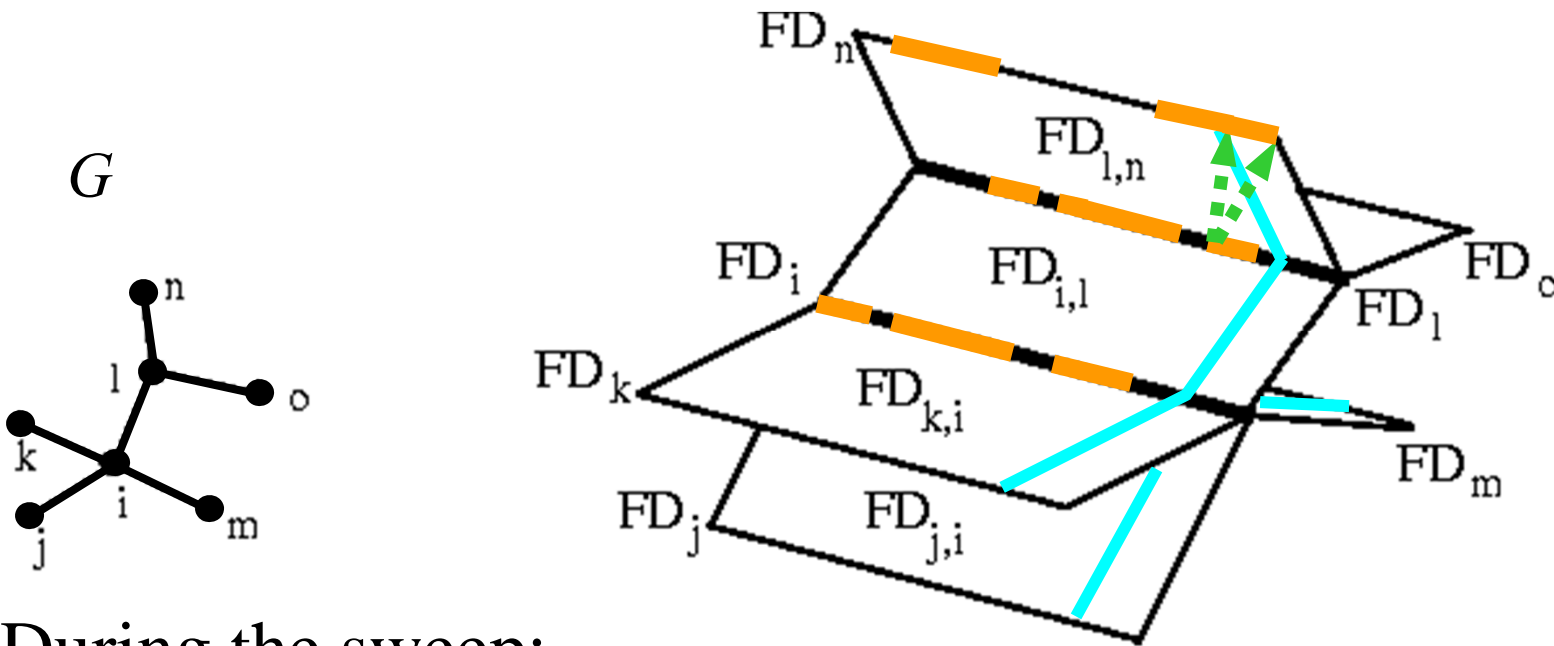
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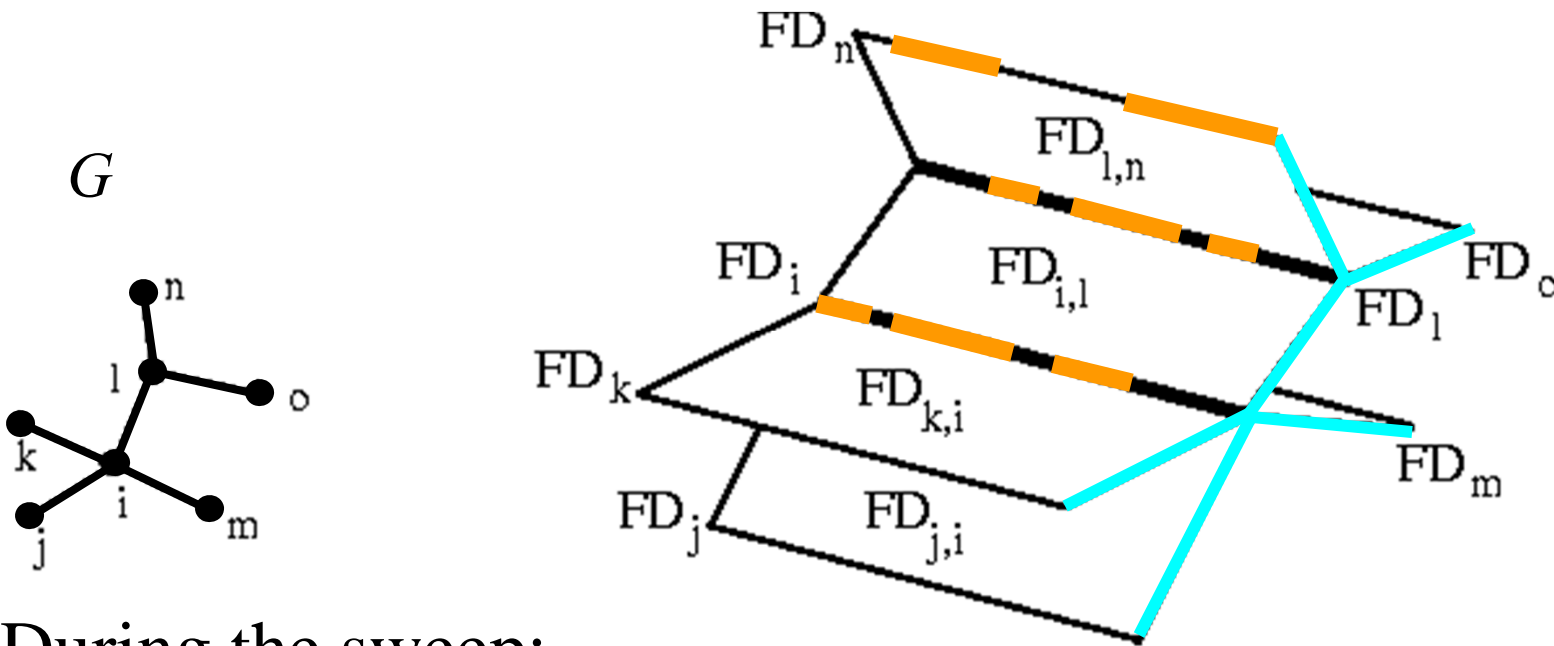
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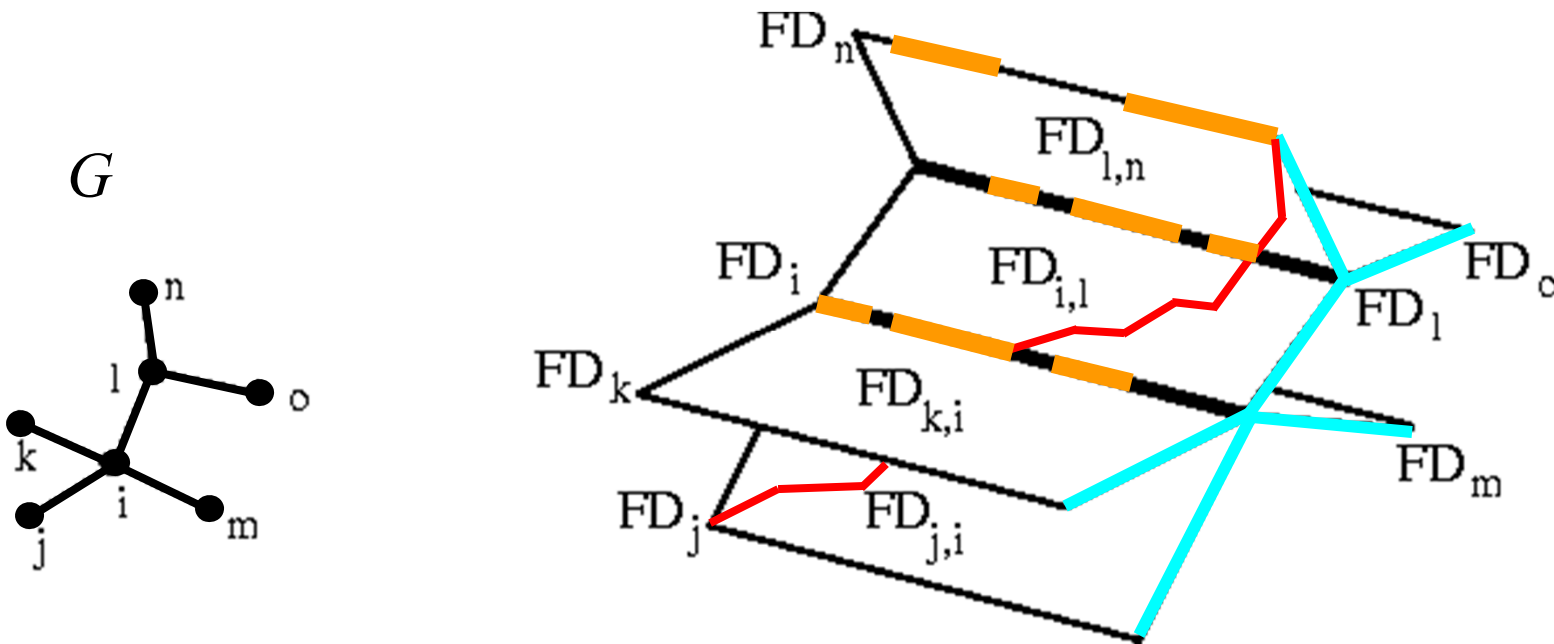
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Backtracking



- After the sweep:
 - Construct a **monotone path via** backtracking

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Map-Matching

- Algorithm for decision problem takes $O(mn \log(mn))$ time and $O(mn)$ space.
- Optimization problem with parametric search: $O(mn \log^2(mn))$ time

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