#### CMPS 3130/6130 Computational Geometry Spring 2015



#### *Motion Planning* Carola Wenk

## **Robot motion planning**

- **Given:** A floor plan (2d polygonal region with obstacles), and a robot (2D simple polygon)
- **Task:** Find a collision-free path from start to end



- Robot *R* is a simple polygon; R = R(0,0)
- Let  $R(x,y) = R + \begin{pmatrix} x \\ y \end{pmatrix}$ *R* translated by  $\vec{x} = \begin{pmatrix} x \\ y \end{pmatrix}$
- Add rotations:  $R(x, y, \varphi)$





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## **Configuration space**

- # parameters = degrees of freedom (DOF)
- Parameter space = "Configuration space" C:
  - 2D translating: configuration space is  $\mathbf{C} = \mathbf{R}^2$
  - 2D translating and rotating: configuration space is  $\mathbf{C} = \mathbf{R}^2 \ge (0, 2\pi)$



- Obstacle P. Its corresponding C-obstacle P'={ $\vec{x} \in C | R(\vec{x}) \cap P \neq \emptyset$ }
- Free space: Placements (= subset of configuration space) where robot does not intersect any obstacle.

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## **Translating a point robot**

- Work space = configuration space
- Compute trapezoidal map of disjoint polygonal obstacles in  $O(n \log n)$  expected time (where n = total # edges), including point location data structure
- Construct road map in trapezoidal map:
  - One vertex on each vertical edge
  - One vertex in center of each trapezoid
  - Edges between center-vertex and edge-vertex of same trapezoid
  - O(n) time and space



# **Translating a point robot**

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  - Edges between center-vertex and edge-vertex of same trapezoid
  - O(n) time and space
- Compute path:
  - Locate trapezoids containing start and end
  - Traverse this road map using DFS or BFS to find path from start to end



**Theorem:** One can preprocess a set of obstacles (with n = total # edges) in  $O(n \log n)$  expected time, such that for any (start/end) query a collision-free path can be computed in O(n) time.

#### Minkowski sums





#### **Extreme points**





# **Configuration space with translations and rotations**







### Shortest path for robot





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## Visibility graph



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