3/19/13

4. Homework Due **Tuesday 4/2/13** before class

1. Worst-Case DT Runtime (5 points)

Give an example that shows that the worst-case runtime of the randomized algorithm to compute the Delaunay triangulation of a set of n points in the plane is $\Omega(n^2)$. (Hint: It might help to play with one of the Delaunay triangulation programs.)

2. Voronoi (10 points)

We saw in class that the Voronoi diagram of a set of points in \mathbb{R}^2 is the projection of the upper envelope of the dual lifted set of planes in \mathbb{R}^3 . What does the projection of the *lower* envelope correspond to? Similarly, what does the projection of the *upper* convex hull of the points lifted to \mathbb{R}^3 correspond to?

Answer these questions by researching on the internet; as usual, cite the source you were using and give an explanation in your own words.

3. Convex Hull of Intersections (10 points)

Let \mathcal{L} be a set of n lines in the plane, no two of which are parallel. Let S be the set of all $O(n^2)$ intersection points between any two lines in \mathcal{L} .

- (a) Give an $O(n \log n)$ time algorithm to compute an axis-parallel rectangle that contains S.
- (b) [**Optional; for extra credit**] Give an $O(n \log n)$ time algorithm that computes CH(S).

(Hint: Your algorithms cannot compute all points in S explicitly. Sort all lines by slope, and prove that it is enough to consider only a certain subset of intersection points.)

4. Linear Separator (10 points)

Let $R = \{r_1, \ldots, r_m\}$ be set of *m* red points, and let $B = \{b_1, \ldots, b_n\}$ be a set of *n* blue points in the plane. A line *l* is called a **linear separator** if all points of *R* lie on one side of *l* and all points of *B* lie on the other side. (You may assume appropriate general position, and may disregard points that lie exactly on the line.)

Use point-line duality to develop an algorithm for this problem which runs in expected linear time. (*Hint: Linear Programming.*)