



# CMPS 2200 – Fall 2012

## *Randomized Algorithms & Quicksort*

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Slides courtesy of Charles Leiserson with additions  
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# Deterministic Algorithms

Runtime for deterministic algorithms with input size  $n$ :

- Best-case runtime
  - Attained by one input of size  $n$
- Worst-case runtime
  - Attained by one input of size  $n$
- Average runtime
  - Averaged **over all possible inputs** of size  $n$

# Deterministic Algorithms: Insertion Sort

Best-case runtime:  $O(n)$ , input  $[1,2,3,\dots,n]$

→ Attained by one input of size  $n$

• Worst-case runtime:  $O(n^2)$ , input  $[n, n-1, \dots, 2, 1]$

→ Attained by one input of size  $n$

• Average runtime :  $O(n^2)$

→ Averaged **over all possible inputs** of size  $n$

• What kind of inputs are there?

• How many inputs are there?

# Average Runtime

- What kind of inputs are there?
  - Do  $[1, 2, \dots, n]$  and  $[5, 6, \dots, n+5]$  cause different behavior of Insertion Sort?
  - No. Therefore it suffices to only consider all permutations of  $[1, 2, \dots, n]$ .
- How many inputs are there?
  - There are  $n!$  different permutations of  $[1, 2, \dots, n]$

# Average Runtime

## Insertion Sort: $n=4$

```

for j=2 to n {
    key = A[j]
    // insert A[j] into sorted sequen
    i=j-1
    while(i>0 && A[i]>key){
        A[i+1]=A[i]
        i--
    }
    A[i+1]=key
}

```

- Inputs:  $4!=24$

[1,2,3,4] <b>0</b>	[4,1,2,3] <b>3</b>	[4,1,3,2] <b>4</b>	[4,3,2,1] <b>6</b>
[2,1,3,4] <b>1</b>	[1,4,2,3] <b>2</b>	[1,4,3,2] <b>3</b>	[3,4,2,1] <b>5</b>
[1,3,2,4] <b>1</b>	[1,2,4,3] <b>1</b>	[1,3,4,2] <b>2</b>	[3,2,4,1] <b>4</b>
[3,1,2,4] <b>2</b>	[4,2,1,3] <b>4</b>	[4,3,1,2] <b>5</b>	[4,2,3,1] <b>5</b>
[3,2,1,4] <b>3</b>	[2,1,4,3] <b>2</b>	[3,4,1,2] <b>4</b>	[2,4,3,1] <b>4</b>
[2,3,1,4] <b>2</b>	[2,4,1,3] <b>3</b>	[3,1,4,2] <b>3</b>	[2,3,4,1] <b>3</b>

- Runtime is proportional to:  $3 + \text{\#times in while loop}$
- Best:  $3+0$ , Worst:  $3+6=9$ , Average:  $3+72/24 = 6$

# Average Runtime: Insertion Sort

- The average runtime averages runtimes over all  $n!$  different input permutations
  - Disadvantage of considering average runtime:
    - There are still worst-case inputs that will have the worst-case runtime
    - Are all inputs really equally likely? That depends on the application
- ⇒ **Better:** Use a randomized algorithm

# Randomized Algorithm: Insertion Sort

- **Randomize the order of the input array:**
  - Either prior to calling insertion sort,
  - or during insertion sort (insert random element)
- This makes the runtime depend on a probabilistic experiment (sequence of numbers obtained from random number generator)
  - ⇒ Runtime is a random variable (maps sequence of random numbers to runtimes)
- **Expected runtime** = expected value of runtime random variable

# Randomized Algorithm: Insertion Sort

- Runtime is independent of input order ([1,2,3,4] may have good or bad runtime, depending on sequence of random numbers)
  - No assumptions need to be made about input distribution
  - No one specific input elicits worst-case behavior
  - The worst case is determined only by the output of a random-number generator.
- ⇒ When possible use expected runtimes of randomized algorithms instead of average case analysis of deterministic algorithms