## CMPS 2200 Introduction to Algorithms - Fall 12

## 4. Homework <br> Due 10/9/12 in the lab

## 1. Best case for quicksort (5 points)

Let "Deterministic Quicksort" be the non-randomized Quicksort which takes the first element as a pivot, using the partition routine that we covered in class on slide 4 of the quicksort slides.
In the best case the pivot always splits the array in half, for all recursive calls of Deterministic Quicksort. Give a sequence of 3 distinct numbers, a sequence of 7 distinct numbers, and a sequence of 15 distinct numbers that cause this best-case behavior. (For the sequence of 15 numbers the first two recursive calls should be on sequences of 7 numbers each, and the next recursive calls on sequences of 3 numbers).

## 2. Decision tree (5 points)

Below is the code for Bubble Sort:

```
void bubbleSort(int A[1..n]){
    for(int i=1; i <= n; i++)
        for(int j=n; j >= i+1; j--)
            if(A[j]<A[j-1])
                swap(A[j],A[j-1]);
}
```

Draw the decision tree for Bubble Sort for an array $A[1 . .3]$ of $n=3$ elements. Annotate the decision tree with comments indicating the part of the algorithm that a comparison belongs to.

## 3. Lower bound for comparison-based searching (5 points)

Consider the problem of searching for a given key in a sorted array of $n$ numbers. Use a decision tree to show a lower bound of $\Omega(\log n)$ for any comparison-based search algorithm. (Hint: The decision tree needs to represent the output of the search algorithm in its leaves. What should be stored in the leaves?)

## 4. Close to the median ( 5 points)

Let $A[1 . . n]$ be an unsorted array of $n$ numbers, and let $k$ be a positive integer $k \leq n$. Assume both $n$ and $k$ are odd. The task is to design an algorithm that outputs the $k$ numbers that are the closest to the median. These are the numbers whose $k$ rank is closest to $n / 2$. For example, if $A=[2,1,7,4,6,5,8]$ and $k=3$ then the output should be the numbers $5,4,6$. Note that the numbers do not necessarily have to be output in sorted order.)
Give pseudocode for such an algorithm that runs in $O(n)$ time (either expected, or worst-case). (Hint: Use the Select algorithm.)

## Practice Problems

## (Not required for homework credit.)

## 1. Quicksort decision tree

The decision tree on the slides in class was for insertion sort of three distinct elements. Draw the decision tree for quicksort of an array $A[1 . .3]$ of three elements. Use the deterministic quicksort and the partition routine described on the slides. Annotate the decision tree with the current state of the array.

## 2. Median computation

Suppose arrays $A$ and $B$ are both sorted and contain $n$ elements each. Give a divide-and-conquer algorithm to find the median of $A \cup B$ in $O(\log n)$ time. (Describe it either in words or as pseudo-code; whatever you prefer). Argue shortly why the runtime is $O(\log n)$.

