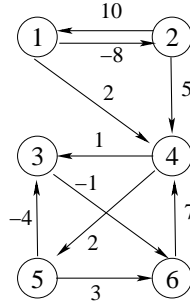


# 10. Homework

Due **12/5/12** in class

## 1. Floyd-Warshall (7 points)

Run the Floyd-Warshall algorithm on the weighed directed graph below. Show the matrix  $c^{(k)}$  that results for each iteration of the outer loop (for  $k=0, \dots, 6$ ).



## 2. To be or not to be ... in NP (8 points)

Which of the problems below are in NP, and which are not? Either justify why the problem is not in NP, or show that it is in NP by sketching an appropriate algorithm and its runtime.

- (a) Given an undirected graph  $G$ . Is  $G$  a tree (i.e., acyclic)?
- (b) Given a positive integer  $a$ , is  $a$  not a prime number (i.e., is  $a$  the product of two integers greater than 1)?
- (c) Given an undirected graph  $G$  with edge-weights. Compute a minimum spanning tree of  $G$ .
- (d) Given an array  $A[1..n]$  of numbers, and a number  $K$ . Is there a subset of numbers in  $A$  that sum up to exactly  $K$ ?

## 3. WEIGHTED CLIQUE (6 points)

Consider the following problem:

**WEIGHTED CLIQUE**

**Given:** An undirected graph  $G$  with positive integer edge weights, and an integer  $S > 0$ .

**Task:** Does  $G$  contain a clique  $C$  with total weight (= sum of edge weights) at least  $S$ ?

Show that WEIGHTED CLIQUE is NP-complete by showing:

- (a) that WEIGHTED CLIQUE is in NP, and
- (b) by giving a polynomial-time reduction from CLIQUE (so, show  $\text{CLIQUE} \leq \text{WEIGHTED CLIQUE}$ )

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4. **Fun with reductions (6 points)**

Suppose  $\Pi_1$  and  $\Pi_2$  are decision problems and  $\Pi_1$  is polynomial time reducible to  $\Pi_2$ , so,  $\Pi_1 \leq \Pi_2$ . Please answer each of the questions below, and justify your answers.

- (a) If  $\Pi_1 \in P$  does this imply that  $\Pi_2 \in P$ ?
- (b) If  $\Pi_2 \in P$  does this imply that  $\Pi_1 \in NP$ ?
- (c) If  $\Pi_1$  is  $NP$ -complete and  $\Pi_2 \in P$ , what does that imply?